

#### Changes and Enhancements

## Acknowledgements

- Larry Woodman Technical
  Director/Kernel Engineering, Mission
  Critical Linux
- Joe Pranevich
- Thomas Wolfgang Burger

## Background

Linux 0.95 in 1991 First release to public via Internet Linux 2.0 Starting to get public notice Linux 2.2 New file systems Redesign of caching Greater scalability Linux 2.4

## Overview

- Linux kernel architecture features
- Linux kernel hardware support features
- File System Enhancements
- Networking Enhancements
- Device Support Enhancements

- ELF and POSIX Foundation
  - More dependent on ELF
  - More POSIX compliant:
    - Clocks and Timers support

- Memory Usage About the same as 2.2 Shared Memory More compliant with Industry standards

  - Introduces a special "shared memory" filesystem

#### 2.2 Page Replacement Problems

- Page eviction
- Simplistic NRU replacement
- Clock algorithm can evict accessed pages
- Sub-optimal reaction to variable load or load spikes after inactivity

#### 2.4 Improvements:

- Finer-grained SMP locking
- Unification of buffer and page caches
- Support for larger memory configurations
- SYSV shared memory code replaced
- Page aging reintroduced
- Active & inactive page lists
- Optimized page flushing
- Controlled background page aging
- Aggressive readahead

#### SMP locking optimizations

- Use of global "kernel\_lock" was minimized.
- More subsystem based spinlock are used.
- More spinlocks embedded in data structures.
- Semaphores used to serialize address space access.
- More of a spinlock hierarchy established.
- Spinlock granularity tradeoffs.

- Increased number CPUs supported
  - Static increase of maximum CPUs to 64.
  - Realistic scalability of up to 8 CPUs.
    - Bus saturation
    - SMP locking
  - Scheduler optimizations speed up selection of threads and context switching.

#### Kernel multi-threading improvements

- Multiple threads can access address space data structures simultaneously.
- Single mem->msem semaphore was replaced with multiple reader/single writer semaphore.
- Reader lock is now acquired for reading per address space data structures.
- Exclusive write lock is acquired when altering per address space data structures.

#### 32 bit UIDs and GIDs

- Increase from 16 to 32 bit UIDs allow up to 4.2 billion users.
- Increase from 16 to 32 bit GIDs allow up to 4.2 billion groups.

#### 64 bit virtual address space

- Architectural limit of the virtual address space was expanded to a full 64 bits.
- IA64 currently implements 51 bits (16 disjoint 47 bit regions)
- Alpha currently implements 43 bits (2 disjoint 42 bit regions)
- S/390 currently implements 42 bits
- Future Alpha is expanded to 48 bits (2 disjoint 47 bit regions)

#### Unified file system cache

- Single pagecache was unified from previous pagecache read/buffermem write functionality
- Eliminates copying buffers from buffermem to pagecache on file read operations.
- Reduces memory consumption by eliminating double buffered copies of file system data.
- Eliminates overhead of searching two levels of data cache.

#### Distributed Interrupts

- Hardware interrupt service routines can be processed simultaneously on all CPUs.
- Software interrupts (softIRQs) can be processed simultaneously on all CPUs.
- SMP spin locks are maintained within device specific data structures.

#### Increased number of threads and tasks

- Default maximum number of tasks/address spaces was increased.
- Default maximum number of threads per task was increased.
- Configuration of both maximums was changed to be runtime tunable via /proc file system.
- Scheduler optimizations minimize overhead of context switching between sibling threads.

- IA64 Port and Architecture Optimizations
  - Support for IA64 processor features:
    - IA64 specific TLB optimizations.
    - Large rotating register file.
    - IA64 SMP specifics.
    - IA64 IO specifics.
  - 64 bit virtual address space.
    - Itanium is actually 51 bits; sixteen 47 bit regions.
  - NUMA support under development.

- Alpha Architecture Optimizations
  - 64 bit virtual address space.
    - EV67 is 43 bits; half user, half kernel.
    - EV7 supports 48 bits; half user, half kernel.
  - **2TB(41 bit) physical address limit.**
  - Highly accurate SMP compatible processor time optimizations.
  - NUMA support under development.

#### S/390 Architecture Optimizations

- 64 bit virtual address space
  - 42 bits used separate address spaces for users & kernel
- 16EB physical address limit.
- Highly accurate SMP compatible processor time optimizations.
- NUMA support under development.

- BIGMEM for IA32 (and other 32 bit systems)
  - 1GB physical memory limitation in the Linux kernel.
  - 4GB physical memory limitation for 32 bit systems.
  - 4GB physical memory optimizations in the Linux kernel.
  - 64GB physical memory using PAE on IA32.

Special instructions for some processors

- Use of processor specific memory transfer instructions for:
  - Intel Pentium
  - AMD
  - Cyrix
  - WinChip

# 2.4 Kernel Hardware Support Features

- NUMA infrastructure
  - Machine independent Non-Uniform Memory Architecture (NUMA) infrastructure.
  - Support for:
    - multiple memory domains
    - processor subsets
    - binding of devices and interrupts to processors
  - Machine dependent NUMA portion under development for multiple architectures.

#### File System size increase

- File system data offset was increased from 31 bits to 44 bits in the VFS layer.
  - Increases file system size to 16TB.
  - Increases individual file size to 16TB.
  - Still need to consider file system overhead...
- Several local file systems have been enhanced to take advantage of larger files.

- VFS layer redesign to use single cache
  - buffermem and pagecache functionality was unified in 2.4
  - VFS layer was changed to use pagecache for generic file read() and write() operations.
  - Eliminated coping between buffermem and pagecache.
  - Saves memory be eliminating multiple copies of buffered file system data.

- RawIO support to bypass file system cache
  - New RawIO interface was added to file systems.
  - This results in:
    - DMA directly to buffer wired in user address space.
    - Bypassing the pagecache.
    - Eliminates coping between pagecache pages and user buffer pages.
    - More efficient for databases.

- Several Journaling File Systems introduced
  - Pending file system updates are continually maintained in a single journal file.
  - The FSCK at reboot time is reduced to replaying the journal.
  - Speeds up reboot FSCK by several orders of magnitude.
  - ext3fs, reiserfs, xfs

- Inclusion of Logical Volume Manager into the Linux kernel
  - Allows file systems to span multiple disks.
  - Dynamic runtime resizing of file systems.
  - More flexible file system device management.
  - Standards compliant.
  - Familiar to users of commercial UNIX.

- Network re-write for optimal performance
  - Redesigned to take advantage of improved multitasking and multithreading.
  - Improvement performance for simultaneous/multiple network interfaces.
  - Distributes networking load much more evenly on SMP systems.
  - Kernel uses wakeup\_one to minimize wasted cycles

# Kernel Networking Enhancements

- Firewall and IP functions placed in kernel
- Network subsystem split:
  - Packet filtering layer
  - Network Address Translation layer
- PPP code rewritten and modularized
- ISDN updated to support many new cards
- PLIP improved
- DECnet & ARCNet protocols supported
- Autodetection of Windows shares based on SMB
- Completely compatible to the letter of IPv4 spec

iptables/netfilter replacement for ipchains

- Linux 2.2 replaced ipfwadm with ipchains.
- Linux 2.4 replaced ipchains with iptables, also known as netfilter.
  - Includes capabilities to construct more sophisticated firewalls.
  - Can be used to implement NAT for supporting masqueraded private networks
  - Compatible with ipfwadm and ipchains command syntax.

#### Kernel based HTTP daemon

- khttpd is a kernel daemon module which serves static web pages.
- Can cooperate with Apache and other web servers to serve dynamic web pages.
- Will result in significant web benchmarking improvement (SpecWeb, etc).

- Fully compatible NFSv3 implementation
  - Fully compatible with version 3 of NFS distributed by Sun Microsystems.
  - Eases the burden of Linux sysadmins who maintain heterogeneous environments.
  - Also compatible with:
    - DECnet
    - ARCnet

- 2.4 Supports:
  - Up to 10 IDE controllers
  - Up to 16 ethernet cards
  - Multiple AIPCs
  - SCSI TCQ (tagged command queuing)
  - RAID devices
  - ATM

#### Buses

- Integrated into the new resource management subsystem
- Plug-N-Play
  - ISA & S/390 device configuration and detection

#### USB

- I2O supported (PCI extension)
- PCMCIA support integrated

#### Framebuffers

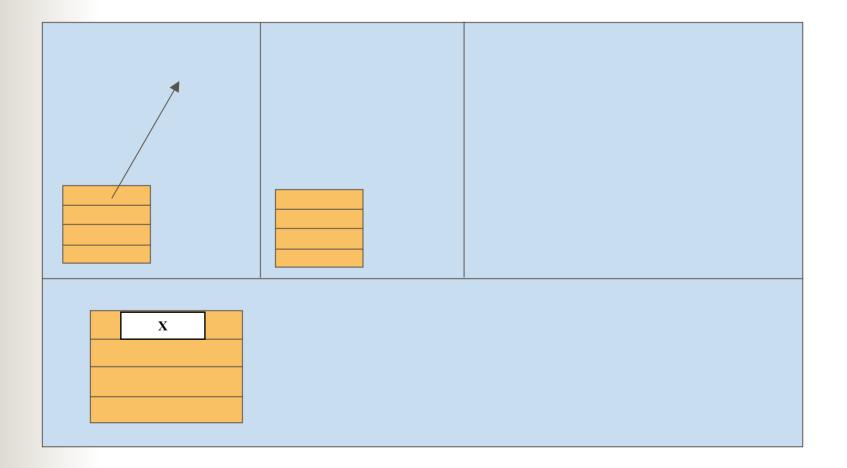
- New drivers and improvements to old
- Support of many more "standard" VGA cards

- Keyboards, Mice, Consoles, and Ports
  - USB support of keyboards and mice
  - Ability to redirect console output to parallel port
  - Serial support has same limitations as 2.2
  - Parallel port support has been overhauled
    - New generic driver
    - DMA support
  - IRDA support
  - Little work done on "WinModems"

- Accessibility
  - Support for speech synthesizer card
- Multimedia
  - No ground breaking changes
  - Updates and new drivers for variety of cards
    - Including full duplex support
  - Ease of configuration enhancements

- S/390 Devices
  - **3270** as console and terminal
  - Tape support
  - Hipersockets
  - z/OS formatted disk (VTOC & DCBs)
  - PAGEX support (VM only)
  - Kernel in NSS (VM only)

# PAGEX/PFAULT Support



# PAGEX/PFAULT Support

- Eliminate overhead of double paging
- Page fault by Linux virtual machine usually puts it in wait state until VM gets page
- PAGEX/PFAULT handshaking allows VM to inform Linux of page request and have it dispatch another process
- When page operation is complete VM signals
  Linux again so it can mark task as dispatchable

# PAGEX/PFAULT Support

#### PAGEX

- PROG 14 interrupt
- 32-bit only
- PFAULT
  - External interrupt (x'2603')
  - **32 & 64-bit systems**
  - z/VM 4.2 required

## What's still needed?

- Greater scalability above 8 processors
- NUMA
- Improved fiber-channel handling (requires an inappropriate amount of hand waving to work)
- >1TB per file system limit
- Poor I/O throughput on x86 class machines with very large amounts of memory
- Basic fail-over is there but not advanced clustering
- Logical volume manager needs more work