

IBM Systems & Technology Group

z/VM Performance Case Studies Session 9166

Please consider sitting near the front.

Bill Bitner VM Performance Evaluation bitnerb@us.ibm.com

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Agenda

- Brief review of collecting performance data.
- Brief review of IBM Performance Support
- Case Studies

Acknowledgments

Thanks to the z/VM Performance Team, particularly Brian Wade, for their contributions to this material.



Collecting Raw Monitor Data

- What is raw monitor data?
- How do I set up to collect it?
- When do I collect it?
- What tools are available to help me collect it?
- How do I package it for transmission?
- How do I study it myself?
- Summary



What is Raw Monitor Data?

- It is unformatted binary data describing system configuration or activity
- Logically, it is a sequence of monitor records
 - Each record comments on some specific aspect of system activity or performance
 - In aggregate they constitute a comprehensive, time-indexed record of system activity

There are three large classes of monitor records

- Configuration records: emitted when monitor starts, these describe system configuration
- Sample records: emitted every so often, these comment on the accumulated activity of an entity (device, user, ...)
- *Event records:* emitted as needed, these comment on some specific phenomenon that just now occurred
- Some records come from the Control Program and comment on its experience in running the system
- Other records come from guests and comment on their experiences in doing whatever it is they do
- We collect this data using an IBM-supplied utility program called MONWRITE
- During the rest of this presentation, we will call this data MONWRITE data

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How Do I Collect MONWRITE Data?

- By Default the z/VM system is set up with DCSS and user ID named MONWRITE
- If somehow skipped, then:
 - You set up a DCSS where CP will buffer the monitor records it emits
 - CP DEFSEG and SAVESEG commands
 - You tell CP which kinds of records to emit, and how often to emit them, and in fact to begin emitting them
 - CP MONITOR command
 - You set up a guest that drains the DCSS to a disk or a tape via the MONWRITE utility
- On some occasions, the default DCSS (named MONDCSS) is too small.
 - See http://www.vm.ibm.com/perf/tips/mondcss.html
- You run the guest
- You archive the resultant files or tapes, so that you have a long-term historical record of system activity and performance



When Do I Collect MONWRITE Data?

- Periodically, collect and archive some data during your peak periods, so that you have a historical record
 - Every Tuesday at 10 AM for an hour
 - Month-end processing
 - Whenever you do that really big thing you do
- When directed by IBM
 - Health check, PMR, crit sit, ESP, whatever



Tool: Running MONWRITE By Hand

- A great idea, assuming you are not running some other performance product
 - If you know what you are doing, you can do both simultaneously
- Create the DCSS to hold the buffered records
- Set up a guest to run our MONWRITE MODULE (collector)
- Issue some CP MONITOR commands to start CP emitting records
 - Enable all samples
 - Enable all events except seeks and scheduler
 - Use a 1-minute sample interval and a 5-second HFS rate
- In your guest, start MONWRITE to collect the data CP's emitting
- To stop collecting, type this: MONWSTOP
- You will end up with one MONWRITE file that you can:
 - Archive for the historical record
 - Analyze yourself with z/VM Performance Toolkit
 - Send to IBM so we can look at it
- There is an option for MONWRITE to close the file at regular times of day and a user exit to process the just-closed file.
- Good references:
 - http://www.vm.ibm.com/perf/tips/collect.html a good cheat sheet
 - <u>z/VM Performance</u>, chapter 9, "Monitoring Performance Using CP Monitor" an excellent writeup of every last detail

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Tool: Brian Wade's LINMON Collector

- At <u>http://www.vm.ibm.com/devpages/bkw/monitor.html</u>
- Based on a modified MONWRITE
- Sets up the DCSS, etc. on its own, using certain assumptions that are probably safe for many systems
- Can be configured to:
 - Account for presence of another performance product
 - Collect for a while then log off
 - Start a new file every so often
 - Keep only the last N files

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Packaging MONWRITE Data For Transmission

- MONWRITE files are binary CMS files, F 4096.
- Just attaching them to an e-mail is NOT recommended.
- The standard z/VM Level 2 process for FTPing files calls for COPYFILE (PACK)
 - This is unnecessary for MONWRITE and VMARC files.
- Always, always, always:
 - Move the files in binary (ASCII is a sure showstopper)
 - Do not use FTP's SITE FIXREC, QUOTE SITE FIXREC, or LOCSITE FIXREC features (error-prone)
- You will probably FTP your data to IBM's receiving server in Boulder, CO
 - Testcase.boulder.ibm.com, cd /toibm/vm
 - Name your file mnemonically and send us a note about it
 - See <u>http://techsupport.services.ibm.com/390/tcprocs.html</u> for additional info on the Testcase process
 - In PMR and/or note be clear as to what is sent and how packaged
- We suggest you use the VMARC file archiver that runs on CMS for very large files or when several files are being sent.
 - Kind of like "zipping" on a PC (compresses, combines)
 - MONWRITE data is very compressible
 - Sometimes you also want to send us a console spool, or some QUERY outputs, or whatever
 - You can package everything into one VMARC archive and just send us that
- There are VMARC instructions near the bottom of http://www.vm.ibm.com/devpages/bkw/monitor.html



Studying MONWRITE Data

- z/VM Performance Toolkit
- Interactively possible, but not so useful
- PERFKIT BATCH command pretty useful
 - Control files tell Perfkit which reports to produce
 - You can then inspect the reports by hand or programmatically
- See <u>z/VM Performance Toolkit Reference</u> for information on how to use PERFKIT BATCH



Other Types of Data Confused with MONWRITE Data

- Asking for "raw VM monitor" data can be confusing.
 - Velocity has their own form of raw monitor data and history files, and even a form that mimics MONWRITE.
 - "VM Monitor" sounds like the "VM:product" often associated with CA products.
 - Performance Toolkit's history, trend, and summary files do not have the same detail.
- Be specific when asking for data.



Monwrite Summary

- MONWRITE data is a comprehensive record of system activity
- It is invaluable in diagnosing performance concerns
- If you ask IBM for performance help, IBM will very likely ask you for MONWRITE data
- Practice collecting and transmitting MONWRITE data when you are not under duress
- Archive your MONWRITE data routinely so that you have a good record of your system's usual behavior
- Learn to use PERFKIT BATCH to generate reports, and get familiar with a few of the basic reports



Performance Support

The typical lines of support:

- 1. Your FTSS (Field Technical Sales Support)
- 2. If FTSS he needs help, he will contact Region Designated Support (RDS)
- 3. If RDS needs help, they will contact Advanced Technical Support (ATS)
- 4. If ATS needs help, they will contact z/VM Development Lab
- You may also have contract for Q&A help



- 1. **z/VM is Doing Fine:** This study shows how to recognize a healthy z/VM system.
- Logical to Physical CPUs: We'll look at measures of processor time in an LPAR environment and how the over commitment of logical to real can affect both LPAR and z/VM.
- **3. Emergency Scan:** We get asked about "emergency scan" from time to time. This case study explains it and tells why seeing it is not always an "emergency" in the literal sense.
- 4. Why Doesn't My System Page Faster?: This system isn't broken, but the customer didn't understand its behavior. The case study illustrates why it's important to know the big picture when trying to discern meaning in measurement data.
- 5. Undersized LPAR: This system is generally short on storage, CPU, and paging. The case study illustrates how to detect it and how to fix it.
- 6. **PAV and MDC:** This customer called in with a performance PMR and we ended up taking an APAR. See how we put the finger on a CP bug using CP monitor data.
- 7. **Paging Difficulties:** This system was grossly under configured for paging. The case study illustrates what we examined and what changes we recommended. It also illustrates what happens when one relieves a constraint: namely, one usually bumps into another one.
- 8. HiperSockets Performance: HiperSockets is thought of as a super high speed connectivity feature. This case study looks at some aspects of that performance and a scenario where it might not be as fast as expected.
- 9. The Grinch that Stole Performance: This case study illustrates how system performance can change when hardware fails. CP Monitor showed where the problem was and pointed the way for a hardware fix.

Case Study: z/VM is Running Fine

8/22/2009



Question from Customer

- Linux on z/VM on 2086
- Java core and heap dumps
- Linux transaction rollback exceptions
- Linux slow response time
- Can you please take a look
- I have MONWRITE data for you
 - ... that's always tempting, so we looked



Basic Things to Check

- Do we have enough CPU
 - FCX225 SYSSUMLG
 - FCX126 LPAR
 - FCX114 USTAT, %CPU
- Do we have enough storage
 - FCX114 USTAT, %PGW
 - FCX113 UPAGE, XSTORE and DASD paging
- Do we have enough SXS storage
 - FCX264 SXSUTIL
- Are we spending too much time in the Control Program
 - FCX225 SYSSUMLG

- Are we paging OK
 - FCX109 DEVICE CPOWNED, paging I/O performance
 - FCX103 STORAGE, page blocking factors
 - FCX113 UPAGE, is XSTORE more active than DASD
- Is I/O performance OK
 - FCX108 DEVICE
 - FCX177 CACHEXT
- Is networking performance OK
 - Find OSD chpids via FCX161
 - FCX215 FCHANNEL
 - FCX240 VSWITCH

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Do We Have Enough CPU: FCX126 and FCX225

FCX126 Run 2008/07/24 12: 46: 56

LPAR Logical Partition Activity

Partition Nr. Upid #Proc Weight Wait-C Cap %Load CPU %Busy %Ovhd %Susp %VMId %LogId Type LMRHA 0 NO 0 . . . 1 2 . 2 41. 7 LMRPROD 2 01 500 NO NO 0 41.9 41.9 41.8 ICF . . . 500 NO 1 52.2 52.2 . 2 52.0 52.1 ICF

FCX225 Run 2008/07/24 12: 46: 56

SYSSUMLG

System Performance Summary by Time

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End Time	Busy	T/V	ture	line	Busy	ged	Activ	/s	msec	Elist	PGOUT	Write	/s
>>Mean>>	46.9	1.02	. 9945	2.0		22	12	30. 1	3.7	. 0	14.0	. 2	. 0
12: 55: 38	31.0	1. 02	. 9946	2.0		22	12	23.9	2.9	. 0	. 0	. 0	. 0
12: 56: 38	41.3	1.02	. 9961	2.0		22	11	27.8	3.3	. 0	. 0	. 0	. 0
12: 57: 38	47.9	1.01	. 9966	2.0		22	11	20. 2	3.1	. 0	. 0	. 0	. 0
12: 58: 38	51.7	1.01	. 9968	2.0		22	11	27.6	3.0	. 0	. 0	. 0	. 0
12: 59: 38	61.5	1.01	. 9968	2.0		22	11	25. 9	3.0	. 0	1.5	. 0	. 0
13: 00: 38	44.7	1.03	. 9944	2.0		22	11	26.8	3.0	. 0	. 0	. 0	. 0
13: 01: 38	51.1	1.02	. 9961	2.0		22	11	38.9	2.5	. 0	. 0	. 0	. 0
13: 02: 38	40. 2	1.02	. 9956	2.0		22	11	27.7	2.4	. 0	. 0	. 0	. 0
13: 03: 38	44.9	1. 02	. 9955	2.0		22	13	32.9	3.2	. 0	. 0	. 0	. 0

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<- This is us

Only one LPAR is using these two engines.

We don't see any particularly high percentages.

Also note T/V is nearly perfect.

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CPU and Storage: FCX114 USTAT

FCX114 Run 2008/07/24 12: 46: 56	USTAT	
	Wait State Analysis by User	
From 2008/07/24 12:54:38		12
To 2008/07/24 14: 04: 38		CP
For 4200 Secs 01: 10: 00	Result of 12345 Run	z/

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										<-S\	/M ar	nd->					<%	Time	spe	nt i	n>
Useri d	%ACT	%RUN	%CPU	%LDG	%PGW	%I OW	%SIM	%TIW	%CFW	%TI	%EL	%DM	%I 0A	%PGA	%LIM	%OTH	QO	Q1	Q2	Q3	E0-3
>System<	22	21	4	0	0	0	0	71	0	0	0	2	1	0	0	0	98	0	0	0	0
LMRDB2P	99	29	4	0	0	0	0	65	0	0	0	0	2	0	0	0	100	0	0	0	0
LMRLN2P	99	8	4	0	0	0	0	86	0	0	0	0	1	0	0	0	100	0	0	0	0
LMRWASP	99	64	3	0	0	0	0	32	0	0	0	0	1	0	0	0	100	0	0	0	0
DTCVSW2	91	0	9	0	0	0	0	91	0	0	0	0	0	0	0	0	100	0	0	0	0
LMRLN1P	87	5	3	0	0	0	0	92	0	0	0	0	0	0	0	0	100	0	0	0	0
PERFSVM	13	0	0	0	0	0	0	11	0	5	0	60	23	0	0	0	40	0	0	0	0
TCPI P	4	0	1	0	0	0	0	99	0	0	0	0	0	0	0	0	100	0	0	0	0

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Do We Have Enough Storage: FCX113 UPAGE

FCX113 Run 2008/07/24 12: 46: 56	UPAGE	
	User Paging Activity and Storage Utilization	
From 2008/07/24 12:54:38		12345
To 2008/07/24 14:04:38		CPU 2
For 4200 Secs 01: 10:00	Result of 12345 Run	z/VM
To 2008/07/24 14: 04: 38		CPU 2

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>System<	. 0	. 0	. 0	. 5	. 0	. 2	. 5	. 0	117135	0	22167	94880	4	27	9374	39
BKRBKUP	. 0	. 0	. 0	. 0	. 0	. 0	. 0	. 0	83	0	0	65	0	0	399	0
BKRCATLG	. 0	. 0	. 0	. 0	. 0	. 0	. 0	. 0	85	0	0	66	0	0	412	0
DI SKACNT	. 0	. 0	. 0	. 0	. 0	. 0	. 0	. 0	1227	0	0	0	0	0	1227	0
DTCVSW1	. 0	. 0	. 0	. 0	. 0	. 1	. 1	. 0	49	0	1	45	0	1	2614	0
DTCVSW2	. 0	. 0	. 0	. 0	. 0	. 0	. 0	. 0	273	0	10	297	8	26	2385	0
EREP	. 0	. 0	. 0	. 0	. 0	. 0	. 0	. 0	78	0	0	0	0	0	1231	0
LMRDB2P	. 0	. 0	. 0	. 0	. 0	. 4	4.7	. 0	1027k	0	269342	758161	4	68	17870	0
LMRLN1P	. 0	. 0	. 0	. 0	. 0	. 0	3.0	. 0	449333	0	116199	333205	0	71	73948	0
LMRLN2P	. 0	. 0	. 2	. 0	. 0	. 2	. 8	. 2	501237	0	73352	427956	4	68	20735	847
LMRWASP	. 0	. 0	. 0	. 0	. 0	2.0	. 2	. 0	593939	0	28710	565308	48	27	74281	0

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Is SXS OK: FCX264 SXSUTIL

 FCX264
 Run 2008/07/24 12: 46: 56
 SXSUTIL

 System Execution Space Utilization

 From 2008/07/24 12: 54: 38

 To
 2008/07/24 14: 04: 38

 For
 4200 Secs 01: 10: 00

 Result of 12345 Run

< System Execution Space Utilization (Pages)													>
	<> Used> P												
	Avail- < CP> < Aliases>											tial	Conti -
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End Time	Pages Queue	s Total	Mapped	Free	0ther	Total	Total	CPLock	0wned	<2GB	>2GB	Queue	Pages
>>Mean>>	524287 51374	2 10545	3063	812	6485	185	28	0	0	7873	2714	186	1851
12: 55: 38	524287 51375	0 10537	3063	813	6477	184	29	0	0	7859	2719	185	1851
12: 56: 38	524287 51375	6 10531	3063	806	6478	184	29	0	0	7860	2719	185	1851
12: 57: 38	524287 51375	4 10533	3063	807	6479	184	29	0	0	7860	2719	185	1851
12: 58: 38	524287 51375	4 10533	3063	806	6480	184	29	0	0	7862	2719	185	1851

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Are We Paging OK: FCX109 DEVICE CPOWNED

FCX109 Run 2008/07/24 12: 46: 56

DEVICE CPOWNED Load and Performance of CP Owned Disks

Page / SPOOL Allocation S	Summary												
PAGE slots available	6609240	SPOC	L slot	s avail	abl e	6	00840						
PAGE slot utilization	O%	SPOC	L slot	utili	zation		15%						
T-Disk cylinders avail.		DUMF	o slots	availa	abl e		0						
T-Disk space utilization	%	DUMF	slot	utiliza	ation		%						
< Device Descr>		<		Ra ⁻	te/s		>	User		Serv	MLOAD	BI ock	%Used
Volume Area	Area Use	d <pa< td=""><td>ige></td><td><sp< td=""><td>></td><td></td><td>SSCH</td><td>Inter</td><td>Queue</td><td>Ti me</td><td>Resp</td><td>Page</td><td>for</td></sp<></td></pa<>	ige>	<sp< td=""><td>></td><td></td><td>SSCH</td><td>Inter</td><td>Queue</td><td>Ti me</td><td>Resp</td><td>Page</td><td>for</td></sp<>	>		SSCH	Inter	Queue	Ti me	Resp	Page	for
Addr Devtyp Serial Type	Extent	6 P-Rds	P-Wrt	S-Rds	S-Wrt	Total	+RSCH	feres	Lngth	/Page	Time	Si ze	Alloc
OD15 3390 LPVPM7 PAGE	1-3338) . C	. 0			. 0	. 0	0	0	13.3	13.3		0
OD16 3390 LPVPM8 PAGE	1-3338) . C	. 0			. 0	. 0	0	0	13.3	13.3		0
OD17 3390 LPVPM9 PAGE	1-3338) . C	. 0			. 0	. 0	0	0	. 8	. 8	15	100
OD18 3390 LPVPMA PAGE	1-3338).C	. 0			. 0	. 0	0	0	11.9	11.9	14	25
OD19 3390 LPVPMB PAGE	1-3338).C	. 0			. 0	. 0	0	0	5. 9	5.9	12	100
0E05 3390 LPVPM1 PAGE	1-3338).C	. 0			. 0	. 0	0	0	5.3	5.3	16	100
OEO6 3390 LPVPM2 PAGE	1-3338	. 0	. 0			. 0	. 0	0	0	2.1	2. 1	19	100
OEO7 3390 LPVPM3 PAGE	1-3338	. 0	. 0			. 0	. 0	0	0	8.3	8.3	12	50
0E08 3390 LPVPM4 PAGE	1-3338).C	. 0			. 0	. 0	0	0	4.9	4.9	16	100
OEO9 3390 LPVPM5 PAGE	1-3338	. 0	. 0			. 0	. 0	0	0	13.3	13.3		0
OEOA 3390 LPVPM6 PAGE	1-3338	. 0	.0			. 0	. 0	0	0	5.2	5.2	14	100



Time In The Control Program: FCX225 SYSSUMLG

- We saw this report already
- T/V ~ 1.02
- T/V = (CP time + guest time) / guest time
- 1.0 is a perfect T/V (CP=0)



I/O Performance: FCX108 DEVICE

DEVI CE

FCX108 Run 2008/07/24 12:46:56

General I/O Device Load and Performance

< Device Descr>	Mdisk F	Pa-	<-Rat	te/s->	<	1	Fime ((msec))	>	Req.	<perc< td=""><td>cent></td><td>SEEK</td><td>Recov ·</td><td><-Thro</td><td>ttle-></td></perc<>	cent>	SEEK	Recov ·	<-Thro	ttle->
Addr Type Label/ID	Links 1	ths	I/0	Avoi d	Pend	Di sc	Conn	Serv	Resp	CUWt	Qued	Busy	READ	Cyl s	SSCH S	Set/s	DI y/s
>> AII DASD <<			. 2	. 0	. 2	2.1	1.4	3.7	3.7	. 0	. 0	0	23	737	0		. 0
0E21 3390 LPLAM1	1	2	4.6	. 0	. 2	4.8	5.7	10.7	10.7	. 0	. 0	5	0	1372	0		
0E00 3390 LPVRM1 CP	50	2	. 4	. 0	. 2	5.1	1.0	6.3	6.3	. 0	. 0	0	0	7	0		
0E02 3390 LPVWM2 CP	43	2	. 1	. 2	. 2	3.7	. 8	4.7	4.7	. 0	. 0	0	0	46	0		
0E22 3390 LPLQM1	1	2	1.5	. 0	. 2	3.2	1.0	4.4	4.4	. 0	. 0	1	0	1910	0		
0E20 3390 LPLDM1	1	2	3.0	. 6	. 2	2.4	. 9	3.5	3.5	. 0	. 0	1	0	884	0		
OD13 3390 >LMRDB2P	0	2	. 7	. 0	. 2	2.6	. 5	3.3	3.3	. 0	. 0	0	83	393	0		
0E23 3390 LPLQM2	1	2	2.7	. 1	. 2	2.3	. 8	3.3	3.3	. 0	. 0	1	0	870	0		
0E19 3390 LPLZM2	1	2	. 0	. 0	. 2	2.5	. 4	3.1	3.1	. 0	. 0	0	50	52	0		
OF2C 3390 >LMRDB2P	0	2	. 5	. 0	. 2	2.0	. 8	3.0	3.0	. 0	. 0	0	3	803	0		
OD12 3390 >LMRDB2P	0	2	. 8	. 0	. 2	2.3	. 4	2.9	2.9	. 0	. 0	0	61	490	0		
ODOF 3390 >LMRDB2P	0	2	2.8	. 0	. 2	2.0	. 6	2.8	2.8	. 0	. 0	1	63	108	0		
OD11 3390 >LMRDB2P	0	2	1.2	. 0	. 2	2.1	. 4	2.7	2.7	. 0	. 0	0	46	399	0		
OD14 3390 >LMRDB2P	0	2	. 5	. 0	. 2	2.0	. 5	2.7	2.7	. 0	. 0	0	81	283	0		
OE26 3390 >LMRDB2P	0	2	. 5	. 0	. 2	1.6	. 8	2.6	2.6	. 0	. 0	0	21	415	0		
OD2D 3390 >LMRDB2P	0	2	1.3	. 0	. 2	. 0	2.2	2.4	2.4	. 0	. 0	0	9	17	0		
0E28 3390 >LMRDB2P	0	2	. 2	. 0	. 2	1.4	. 8	2.4	2.4	. 0	. 0	0	71	697	0		
0D00 3390 RM1LPV	0	2	. 0	. 0	. 2	. 0	2.1	2.3	2.3	. 0	. 0	0			0		
0D10 3390 >LMRDB2P	0	2	2.2	. 0	. 2	1.6	. 5	2.3	2.3	. 0	. 0	1	17	420	0		

OE21 LPLAM1 and OE00 LPVRM1 a little slow, but I/O rates are so low... worth studying the workload.

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Networking Performance: FCX161 and FCX215

FCX161 Run 2008/07/24 12: 46: 56	LCHANNEL
	Channel Load and Channel Busy Distribution
From 2008/07/24 12:54:38	
To 2008/07/24 14:04:38	
For 4200 Secs 01: 10: 00	Result of 12345 Run

CHPID Chan-Group		<%Bus	usy> < Channel %Busy Distribution 12:54:38-14:04:38									>		
(Hex)	Descr	Qual Shrd	Cur /	Ave	0-10	11-20	21-30	31-40	41-50	51-60	61-70	71-80	81-90	91-100
11	OSD	00 Yes	15	10	77	23	0	0	0	0	0	0	0	0
00	OSD	00 Yes	0	0	100	0	0	0	0	0	0	0	0	0
01	OSD	00 Yes	0	0	100	0	0	0	0	0	0	0	0	0
10	OSD	00 Yes	0	0	100	0	0	0	0	0	0	0	0	0

FCX215 Run 2008/07/24 12: 46: 56

FCHANNEL

FICON Channel Load

<----- FICON Channel Utilization % ----->

		< To	otal fo	or Syst	em>	<-Own Partition>			<total< th=""><th></th></total<>		
Channel		Bus	Work	<data< td=""><td>Units></td><td>Work</td><td><data< td=""><td>Uni ts></td><td><transfe< td=""><td>r Rate-></td><td></td></transfe<></td></data<></td></data<>	Units>	Work	<data< td=""><td>Uni ts></td><td><transfe< td=""><td>r Rate-></td><td></td></transfe<></td></data<>	Uni ts>	<transfe< td=""><td>r Rate-></td><td></td></transfe<>	r Rate->	
Path		Cycl e	Uni ts	Write	Read	Uni ts	Write	Read	<- (Byte	s/s)>	
I D	Shrd	T_BCy	T_WUn	T_DUW	T_DUR	L_WUn	L_DUW	L_DUR	Write/s	Read/s	
11	Yes	0	10	0	0	10	0	0	513697	127155	< 1 MB/sec altogether
50	Yes	0	0	0	0	0	0	0	259435	25383	
40	Yes	0	0	0	0	0	0	0	250804	25338	

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Networking Performance: FCX240 VSWITCH

ТСН
TCH Activity
ult of 12345 Run
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·		•	•	•	•	•	•	•	•	•	•	· ·
	QT	ime	< 0u	tbound/s	>	< Ir	nbound/s	s>	< S	i gnal s	>	< Intrpts >
	S	Out	Bytes 🗸	<packet< td=""><td>S></td><td>Bytes</td><td><pack< td=""><td>kets></td><td>< is</td><td>sued/s</td><td>></td><td></td></pack<></td></packet<>	S>	Bytes	<pack< td=""><td>kets></td><td>< is</td><td>sued/s</td><td>></td><td></td></pack<>	kets>	< is	sued/s	>	
Addr Name Co	ontrolr V	Sec	T_Byte ⁻	T_Pack T_	Di sc	R_Byte	R_Pack	R_Disc	Write	Read	Sync	Rcv/s Pro/s
>> System	<< 8	300	502784	463.6	. 0	116135	351.0	. 0	165.0	. 0	. 0	244.4 238.8
02F2 D	TCVSW2 8	300	502784	463.6	. 0	116135	351.0	. 0	165.0	. 0	. 0	244.4 238.8



Summary

- There doesn't seem to be anything wrong with this z/VM
- It's worth looking inside the Linux guests
- The questioner directed to Linux support

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Case Study: Logical to Physical

8/22/2009



Logical to Physical Processor Ratios

- As the number of partitions and their size increases, questions continue to arise as to how to configure z/VM systems
- This case study illustrates some of the factors and information that can be examined
- More complex scenarios would include mixed engine environments



Configuration

- **2097-401**
- 18 Physical Processors
 - 1 CP
 - 17 IFLs
 - 3 Dedicated
- 11 Partitions
 - -5 Active Shared: 3+3+13+13+2 = 34 logicals IFLs
- Ratio of Non-dedicated Logical to Physical CPUs: 2.4



Partition Configs

<Partition->

Name	Nr.	Upi d	#Proc	Weight	Wait-C	Сар
DRLPAR	1		0	0	NO	NO
A5Q1	2		0	0	NO	NO
A5Q2	3		0	0	NO	NO
A5Q3	4		0	0	NO	NO
A5Q4	5		0	0	NO	NO
A5T	6	15	2	4	NO	NO
A5X	7	14	3	DED	YES	NO
LPAR1	8	01	3	2	NO	NO
LPAR2	9	02	13	46	NO	NO
LPAR3	10	03	3	2	NO	NO
LPAR4	11	04	13	46	NO	NO

- 14 undedicated IFLs
- LPAR2 weight equates to 6.44 IFLs
 - **6.44** << 13

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Looking at Processor Time – CPU FCX100

PROC	TYPE	%CPU	%CP	%EMU	%WT	%SYS	%SP	%SIC	%LOGLD		
P00	I FL	46	6	40	54	4	2	77	65	%CPU: total cycles consumed in z/VM.	
P12	I FL	46	5	41	54	3	3	76	65	%CP: total cycles in z/VM	
P11	I FL	46	5	41	54	3	3	76	65	control program	
P01	I FL	46	5	41	54	3	3	76	65	%EMU: total cycles inside	
P02	I FL	46	5	41	54	3	2	77	65	z/VM guests	
P03	I FL	46	5	40	54	4	2	76	65	%SYS: total cycles in CP not associated with a guest	
P04	I FL	46	5	41	54	3	3	76	65	(subset of %CP)	
P05	I FL	46	5	41	54	3	2	76	65	%SP: wall clock time in formal spin locks	
P06	I FL	46	5	40	54	4	2	76	65	•	
P07	I FL	46	5	41	54	3	3	77	65	%LOGLD: pct busy time of time z/VM timers are running	
P08	I FL	46	5	41	54	3	3	76	65		
P09	I FL	46	5	41	54	3	3	76	65		
P10	I FL	46	5	41	54	3	3	77	65		

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<u> </u>		

Looking at Processor Time – LPAR FCX126

%Load	CPU	%Busy	%0∨hd	%Susp	%VMI d	%LogI d	Туре	
34.0	0	47.2	1.4	29.8	45.6	64.9	IFL	%Busy: Total CPU
	1	47.1	1.3	29.8	45.6	64.9	I FL	%0vhd: LPAR Mgmt Time for LCPU
	2	47.1	1.3	29.8	45.5	64.8	I FL	%VMId: %CPU from FCX100
	3	47.2	1.4	29.9	45.6	64.9	I FL	%Susp: <u>1</u> 00% - total of
	4	47.1	1.3	29.8	45.5	64.9	IFL	z/VM Timers
	5	47.1	1.4	29. 9	45.5	64.8	I FL	%Logld: %Logld from FCX100
	6	47.1	1.3	29.8	45.5	64.8	I FL	
	7	47.1	1. 2	29.8	45.6	64.9	IFL	
	8	47.1	1. 2	29.8	45.7	65.0	IFL	
	9	47.1	1.3	29.8	45.6	65.0	IFL	
	10	47.1	1. 2	29.8	45.7	65.0	I FL	
	11	47.0	1. 2	29.7	45.6	64.8	I FL	
	12	47.1	1.2	29.8	45.6	65.0	IFL	



LPAR Mgmt Time (Overhead)

- Solution *** Solution **** Solution *** S
- General LPAR overhead also reported, not associated with a given partition.
- Mgmt time can be influenced by activity and requests from within the partitions



LPAR Suspend Time

- An approximation of when z/VM partition is removed from running for either:
 - Being capped
 - Running other partitions
 - z/VM giving up time via diagnoses while waiting on locks

Another side effect of high suspend time

- z/VM User State Sampling could be skewed



Reconfigure the Logical Processor Counts

Phys	Ded.	LCPUs	Log: Phy	%LPBUSY	%LPOVHD	%NCOVHD	%BUSY	%SUSP
17	3	34	2.4	1249	41	26	1316	29.8%
17	3	24	1.7	851	18	18	887	3.5%

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FCX265 LOCKLOG

	<	Before-	>	<	After -	>
Interval	Locks A	Average	Pct	Locks A	Average	Pct
End Time LockName	/sec	usec	Spi n	/sec	usec	Spi n
>>Mean>> SRMATDLK	563.8	71.78	. 311	474.8	29.76	. 157
>>Mean>> RSAAVCLK	. 0	458.6	. 000	. 0	1.306	. 000
>>Mean>> RSA2GCLK	. 0	187.3	. 000	. 1	6. 128	. 000
>>Mean>> BUTDLKEY	. 0	145.0	. 000	. 0	. 243	. 000
>>Mean>> HCPTMFLK	. 0	. 000	. 000	. 0	. 000	. 000
>>Mean>> RSA2GLCK	6.6	63.55	. 003	16.8	8.880	. 002
>>Mean>> HCPRCCSL	. 0	. 000	. 000	. 0	. 000	. 000
>>Mean>> RSASXQLK	2.9	61.99	. 001	3.1	11.17	. 000
>>Mean>> HCPPGDML	. 5	174.9	. 001	. 7	26.71	. 000
>>Mean>> NSUIMGLK	. 0	. 000	. 000	. 0	. 000	. 000
>>Mean>> FSDVMLK	4.3	39.73	. 001	6.8	14.62	. 001
>>Mean>> HCPPGDPL	1.5	190.9	. 002	1.7	81.73	. 002
>>Mean>> SRMALOCK	. 0	. 000	. 000	. 0	. 000	. 000
>>Mean>> HCPTRQLK	434.5	51.29	. 171	306.0	3.439	. 012
>>Mean>> SRMSLOCK	3062	89.98	2. 119	2193	20. 15	. 491



Summary

- Various rules of thumbs for Logical to Physical
- Starting points
- Look at data
- Suspend time is helpful but has multiple causes

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Case Study: Emergency Scan

8/22/2009

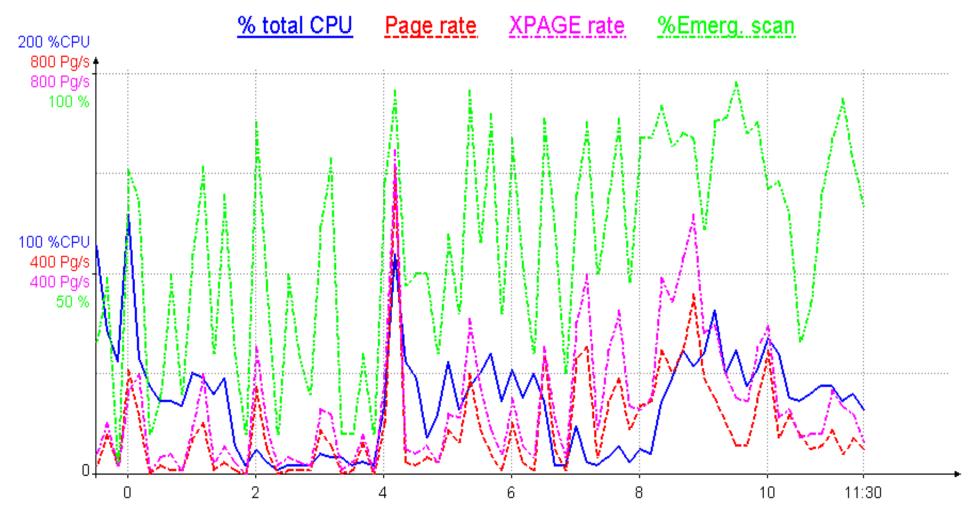


Question from Customer

- My system seems to have a high percentage of emergency scan
- Application performance doesn't seem bothered
- Should I be worried?



Graph from Customer



Source data: Storage

8/22/2009



Finding a Memory Frame

Pass 1: tries to be friendly to dispatched users

- Unreferenced shared address space pages
- Long-term-dormant users
- Eligible-list users
- Dispatch-list users' unreferenced pages down to WSS
- Pass 2: a little more aggressive... like pass 1 except:
 - Avoids shared address spaces
 - Will take from dispatch-list users down to their SET RESERVE
- Emergency scan: anything we can find
- Bit of a misnomer
- Want to know more? Read the prologue of HCPALD

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Is Emergency Scan A Sign of Duress?

Not alone, no.

Evaluate some other things too.

- Are free frame lists routinely zero? (FCX254 AVAILLOG)
- Is system T/V high? (FCX225 SYSSUMLG)
- Are we spinning significantly on any locks? (FCX265 LOCKLOG)
- Does USTAT show users in page wait? (FCX114 USTAT)
- Is an eligible list forming? (FCX100 CPU)
- Are MDC hits satisfactory? (FCX103 STORAGE, FCX108 DEVICE)
- Do you have plenty of SXS space? (FCX264 SXSUTIL)
- Is DASD page rate > XSTORE page rate? (FCX143 PAGELOG)
- Are there queues at paging DASD? (FCX109 DEVICE CPOWNED)
- Is paging MLOAD OK? (FCX109 DEVICE CPOWNED)
- Is paging blocking factor OK? (FCX103 STORAGE)
- Is paging space too full? (FCX109 DEVICE CPOWNED)
- Does application performance seem OK? (you tell me)



Storage Management and VDISKs

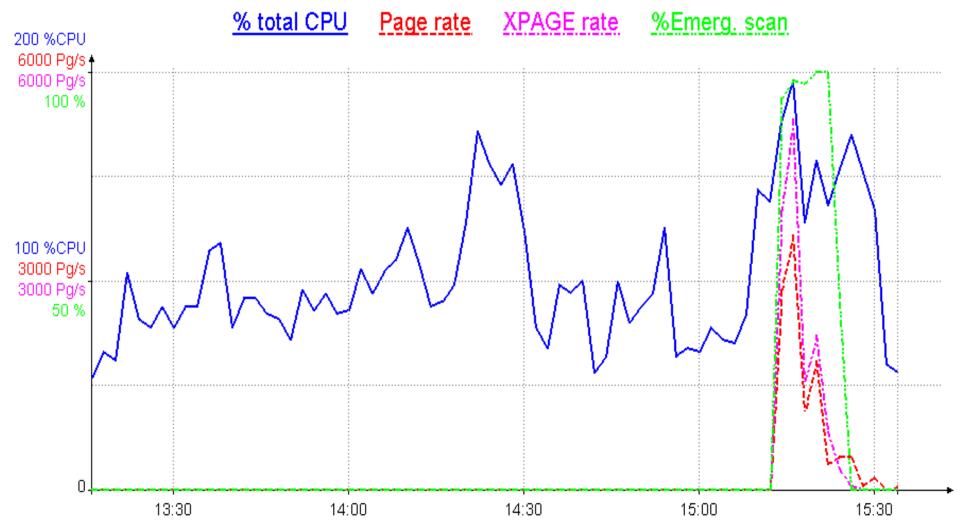
- Referenced VDISK pages are avoided in Pass 1
- This customer realized he had a lot of VDISK for Linux swap space
- If those VDISK pages are used often, they will tend to stick and be ejectable by only emergency scan
- Hmm, customer tried an experiment...

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Customer Removed His VDISKs



Source data: Storage

8/22/2009



Summary

- Try to look at system as a whole
- Whether applications seem debilitated is the best indicator of whether the system is suffering



Case Study: Why Doesn't My System Page Faster



Question from Customer

- "z/VM pages extremely slowly"
- Inactive Linux guest is paged in at only about 1000 pages per second
- 12 3390-9 paging packs, 2 LCUs, with 6 FICON chpids
- During busy periods of running 30 guests, he sees 6000 pages per second
- Customer thinks this single guest should page in much faster
- He devised a 300 MB thrasher that reproduced the behavior
- He sent us lots of charts and graphs
- We asked for MONWRITE data

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Customer Sent MONWRITE Data

- User LIN102 is running the 300 MB thrasher
- It touched 64,000 pages in 61 seconds (1049/sec)
- The interesting time period is 15:12:30 to 15:13:20
- He used MONITOR SAMPLE 10 SEC (brilliant!)
- Ran his data through PERFKIT BATCH
- Looked at some interesting reports for that period



User Configuration

FCX226 UCONF – user configuration report													
					<	- Share -	>				No	Stor	
		Virt	Mach	Stor		%	Max.	Max.	Max.	QUI CK	MDC	Si ze	Reserved
Useri d	SVM	CPUs	Mode	Mode	Relative	Absol ute	Val ue/%	Share	Limit	DSP	Fai r	(MB)	Pages
LI N102	No	1	EME	V=V	100			•••		No	No	768M	0

Virtual uniprocessor with one process (thread) running the memory initializer.

Implications:

- 1. Memory initializer will touch pages serially.
- 2. Page faults will happen serially.



Activity on Paging DASD

FCX108 INTERIM DEVICE 15: 12: 40 to 15: 12: 51

< Device D	escr>	Mdisk	Pa-	<-Rat	te/s->	<		Ti me	(msec)		>	Req.	<perc< th=""><th>cent></th><th>SEEK</th></perc<>	cent>	SEEK
Addr Type	Label /I D	Li nks	ths	I/0	Avoi d	Pend	Di sc	Conn	Serv	Resp	CUWt	Qued	Busy	READ	Cyl s
9F11 3390	VSPPG8 CP	0	6	25.5	. 0	. 2	. 0	3.9	4.1	4.1	. 0	. 0	10	0	131
A062 3390	VSPPG5 CP	0	6	25.0	. 0	. 2	. 0	3.3	3.5	3.5	. 0	. 0	9	100	2580
A02D 3390	VSPPG3 CP	0	6	27.4	. 0	. 2	. 1	3. 1	3.4	3.4	. 0	. 0	9	100	505
9F41 3390	VSPPGB CP	0	6	29.8	. 0	. 2	. 0	3.0	3. 2	3.2	. 0	. 0	10	100	753
A03D 3390	VSPPG2 CP	0	6	35.4	. 0	. 2	. 0	2.9	3.1	3.1	. 0	. 0	11	100	832
9F01 3390	VSPPG7 CP	0	6	38.0	. 0	. 2	. 0	2.8	3.0	3.0	. 0	. 0	11	0	1174
9F5A 3390	VSAPAG CP	0	6	40.9	. 0	. 2	. 0	2.7	2.9	2.9	. 0	. 0	12	100	33
A05D 3390	VSPPG6 CP	0	6	38.9	. 0	. 2	. 0	2.7	2.9	2.9	. 0	. 0	11	100	1446
A01B 3390	VSPPG4 CP	0	6	32.3	. 0	. 2	. 0	2.5	2.7	2.7	. 0	. 0	9	100	2670
9F21 3390	VSPPG9 CP	0	6	45.6	. 0	. 2	. 0	2. 2	2.4	2.4	. 0	. 0	11	0	0
9F51 3390	VSPPGC CP	0	6	48.5	. 0	. 2	. 0	2.2	2.4	2.4	. 0	. 0	12	100	2971
		TOTAL	- 3	387.3									115		

El even paging devices:

- 1. Each in the neighborhood of 10% busy, all reads
- 2. Each showing response time of about 3.1 msec

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Who Else is Doing Paging Activity?

FCX113 UF	PAGE								
	Data	<	F	^p agi ng A	cti vi ty	y/s		>	
	Spaces	<page< td=""><td>Rate></td><td>Page</td><td><pag< td=""><td>ge Mig</td><td>gratio</td><td>on></td><td>Nr of</td></pag<></td></page<>	Rate>	Page	<pag< td=""><td>ge Mig</td><td>gratio</td><td>on></td><td>Nr of</td></pag<>	ge Mig	gratio	on>	Nr of
Useri d	Owned	Reads	Write	Steal s	>2GB>	X>MS	MS>X	X>DS	Users
>System<	. 0	2.3	1.6	7.2	. 0	4.6	6.3	1.7	44
User Data LIN102		75.8	. 0	. 0	. 0	35.2	4.5	. 0	

44 * 2.3 = 101 pages read/sec al together. LIN102 accounts for 76% of this, 76 pages read/sec.



What We Know So Far

- Each paging I/O takes about 3.1 msec
- One single-threaded application in one guest is responsible for most of the paging I/Os
- This means we should see about (1000/3.1) = 323 SSCH ops for paging per second
- We actually saw 387/sec, but remember other guests are paging slightly
- Because one single-threaded guest is responsible for most of the paging I/O, the paging device utilizations should add to about 100%
- They actually add to 115%, but remember other guests are paging slightly

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What Did We Tell The Customer?

- LIN102's page reading speed is limited by its single-threaded nature and the speed of the paging DASD.
- Your system pages at higher rates when 30 guests are running because with multiple guests you can generate concurrent page reads. You have multiple paging exposures too and so you can parallelize paging I/O.
- Your 11 paging exposures look like they could support (1100%/115%) = 9.5 such thrashers concurrently.
- But from FCX109 DEVICE CPOWNED, we see your page space is about 15% full so I wouldn't try more than four of them at once.

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Something Interesting About LIN102

FCX163 Run 2008/05/19 12: 18: 57	UPAGELOG LI N102
	User Paging Activit
From 2008/05/15 15:10:10	
To 2008/05/15 15: 15: 50	
For 340 Secs 00: 05: 40	

Page Data Log for User LIN102

	Data	< Paging Activity/s>									
Interval	Spaces	<page< td=""><td>Rate></td><td><paę< td=""><td colspan="5">age Migration></td></paę<></td></page<>	Rate>	<paę< td=""><td colspan="5">age Migration></td></paę<>	age Migration>						
End Time	Owned	Reads	Write	Steal s	>2GB>	X>MS	MS>X	X>DS			
15: 12: 40	0	437	. 0	. 0	. 0	116	4.2	. 0			
15: 12: 50	0	534	. 0	. 0	. 0	167	. 6	. 0			
15: 13: 00	0	440	. 0	. 0	. 0	342	37.7	. 0			
15: 13: 10	0	313	. 0	. 0	. 0	288	. 2	. 0			
15: 13: 20	0	473	. 0	. 0	. 0	246	3.4	. 0			
Avg		439				232					

Thrasher touched 1049/sec al together.

- 1. 439/sec read from disk
- 2. 232/sec read from XSTORE
- 3. 378/sec resident

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A Note on User States

FCX164 Run 2008/05/19 12: 18: 57

From2008/05/1515: 10: 10To2008/05/1515: 15: 50For340Secs00: 05: 40

USTATLOG LIN102 User Wait States

Wait State Data Log for User LIN102

Interval									
End Time	%ACT	%RUN	%CPU	%LDG	%PGW	%I OW	%SIM	%TIW	%CF
15: 12: 30	100	0	0	0	100	0	0	0	
15: 12: 40	100	0	0	0	100	0	0	0	
15: 12: 50	100	0	0	0	100	0	0	0	
15: 13: 00	100	0	0	0	100	0	0	0	
15: 13: 10	100	0	0	0	100	0	0	0	
15: 13: 20	100	0	0	0	100	0	0	0	

Customer said this means LIN102 "is in page wait 100% of the time".

This is not correct.

It means 100% of the times we looked, LIN102 was in a page wait.

We looked only once every two seconds (FCX149 MONSET).

After all, LIN102 was also touching pages.



Summary

- Customer became absorbed with z/VM measurements but forgot what his workload does
- Knowledge of the workload's behavior is crucial in understanding why the system performs the way it does
- Customer was very good at collecting raw monitor data appropriate for the diagnosis task at hand
- Fun question that was not too difficult to answer

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Case Study: Undersized LPAR

8/22/2009



Question from Customer

Why do my workloads run so slowly?



Customer's Configuration

System model	2094-606
Processors	2 IFL
SYSGEN storage	19968 (19.5 GB)
XSTORE	4096 (4 GB)
Page slots	24641k (94 GB)
Paging devices	25
Logged-on virtual	68 GB



What We Saw in Customer's Data

- Long queues and long response times for paging devices
- Possibility for processor contention during peak hours



BEFORE: Customer's FCX109 DEVICE CPOWNED

FCX109 Run 2008/05/02 12: 13: 56		E CPOWNED and Performance	of CP Ow	vned Disks			
From 2008/04/30 09:50:08							
To 2008/04/30 23: 53: 33							CPU 209
For 50604 Secs 14:03:24							z/VM
Page / SPOOL Allocation Summar	У						
PAGE slots available 246	941k SF	POOL slots avail	abl e	600840			
PAGE slot utilization	37% SF	POOL slot utiliz	zation	31%			
T-Disk cylinders avail	Dl	UMP slots availa	abl e	0			
T-Disk space utilization	% DU	UMP slot utiliza	ation	%			
· ·							
< Device Descr>	<	Rat	te/s	>	User	Serv MLOAD	Block %Used
Volume Area Area	Used <	-Page> <spo< td=""><td>ol></td><td>SSCH</td><td>Inter Queue</td><td>Time Resp</td><td>Page for</td></spo<>	ol>	SSCH	Inter Queue	Time Resp	Page for
Addr Devtyp Serial Type Exte	ent % P-F	Rds P-Wrt S-Rds	S-Wrt To	otal +RSCH	feres Lngth	/Page Time	Size Alloc
7904 3390 520PG7 PAGE C	-3338 57 2 ⁻	1.7 16.9	3	38.6 17.7	1 3.24	6.4 29.0	3 78
790D 3390 520SPL ?????							
SPOOL 1	-3338 31	.0.0.0	. 0	.0.0	1 0	3.7 3.7	100
791B 3390 520PG4 PAGE C	-3338 59 23	3.0 18.0	4	11.0 19.8	1 2.10	5.7 12.9	2 83
7921 3390 52PG14 PAGE C	-3338 59 22	2.7 17.7	4	40.4 19.3	1 3.03	6.0 28.4	2 82
7922 3390 52PG15 PAGE 0	-3338 60 23	3.1 18.0	4	11.1 19.8	1 3.77	5.8 27.6	2 84
792C 3390 520PGA PAGE 8	03060 22 25	5.7 19.6	4	45.4 18.6	1 2.93	5.2 15.7	3 90
792D 3390 520PGC PAGE 8	03060 22 25	5.5 19.7	4	15.2 18.4	1 2.38	4.9 15.5	3 90
7934 3390 520PG8 PAGE 8	03060 22 25	5.5 19.6	4	45.1 18.4	1 2.78	5.0 17.1	3 89
79FC 3390 52PG10 PAGE C	-3338 56 2 ⁻	1.5 16.8	3	38.3 17.6	1 3.29	6.6 26.5	3 78



BEFORE: Customer's FCX225 SYSSUMLG

FCX225 Run 2008/05/02 12: 13: 56	SYSSUMLG	
	System Performance Summary by Time	
From 2008/04/30 09:50:08		VMLNX1
To 2008/04/30 23: 53: 33		CPU 2094-0
For 50604 Secs 14:03:24		z/VM V.

	<	CPU		>	<vec></vec>	<use< th=""><th>rs></th><th>< /</th><th>′0></th><th><stg></stg></th><th><-Pagi</th><th>ing></th><th><spl></spl></th><th><</th><th> UP</th><th>+MP Tra</th><th>nsactio</th><th>ons</th></use<>	rs>	< /	′0>	<stg></stg>	<-Pagi	ing>	<spl></spl>	<	UP	+MP Tra	nsactio	ons
		<rat< td=""><td>i 0></td><td></td><td></td><td></td><td></td><td>SSCH</td><td>DASD</td><td>Users</td><td><-Rate</td><td>e/s></td><td></td><td><-Resp</td><td>onse</td><td>Γime-></td><td><-Trans</td><td>sactio</td></rat<>	i 0>					SSCH	DASD	Users	<-Rate	e/s>		<-Resp	onse	Γime->	<-Trans	sactio
Interval	Pct		Cap-	0n-	Pct	Log-		+RSCH	Resp	in	PGI N+	Read+	Pages		Non-	Qui ck		Non-
End Time	Busy	T/V	ture	line	Busy	ged	Activ	/s	msec	Elist	PGOUT	Write	/s	Tri v	Triv	Di sp	Tri v	Tri v
>>Mean>>	46.7	1. 25	. 8177	2.0		31	26	566.8	11.8	. 0	1680	1046	. 0	. 947	47.16	1. 302	1.15	. 45
10: 00: 08	71.3	1.45	. 7064	2.0		34	30	1071	14.0	. 0	3215	2094	. 0	1.057	50. <mark>79</mark>	1. 558	. 91	. 51
10: 10: 08	68.0	1. 38	. 7363	2.0		34	30	968.9	13.4	. 0	2531	1888	. 2	. 754	48. 21	1. 647	1. 27	. 56
10: 20: 08	71.6	1. 25	. 8138	2.0		30	26	1034	14.3	. 0	2235	1919	. 1	. 966	25.24	1. 382	1.09	. 44
10: 30: 08	37.4	1.17	. 8732	2.0		30	25	418.9	7.7	. 0	957.7	420.8	. 1	. 902	3. 662	1.379	1.09	. 49
10: 40: 08	28.6	1. 28	. 7995	2.0		30	25	332.1	7.8	. 0	1186	504.1	. 0	. 824	11. 15	. 996	1. 28	. 41
10: 50: 08	28.0	1. 25	. 8183	2.0		30	25	328. 1	8. 1	. 0	1023	548.3	. 0	. 982	3.409	1. 261	1.31	. 35
11: 00: 08	32.6	1. 21	. 8433	2.0		30	27	439.1	9.3	. 0	1381	794.2	. 0	. 931	19. 85	1.360	1. 28	. 34
11: 10: 08	31.6	1. 22	. 8385	2.0		30	24	400.4	8.8	. 0	1473	716.6	. 1	. 856	5. 452	1.460	1.37	. 40
11: 20: 08	39.4	1. 26	. 8079	2.0		30	25	571.9	11. 2	. 0	1761	1092	. 0	. 710	9. 936	1.313	1.33	. 38
11: 30: 08	37.0	1. 18	. 8603	2.0		30	24	405.5	9.2	. 0	1324	687.0	. 1	. 766	113. 9	1. 084	1.34	. 54
11: 40: 08	54.7	1. 16	. 8757	2.0		30	25	430.4	9.5	. 0	1395	774.0	. 0	. 849	6. 610	. 907	1.19	. 46
11: 50: 08	57.1	1. 22	. 8318	2.0		30	24	626.6	11.5	. 0	2343	1202	. 0	. 868	13.42	1. 311	1.23	. 45
12: 00: 08	47.7	1.37	. 7433	2.0		30	26	662.0	12. 2	. 0	2185	1238	. 0	. 836	481.5	. 948	1. 28	. 39



BEFORE: Customer's FCX126 LPAR

 FCX126
 Run 2008/05/02 12: 13: 56
 LPAR

 Logical
 Partition Activity

 From 2008/04/30
 09: 50: 08

 To
 2008/04/30
 23: 53: 33

 For
 50604
 Secs 14: 03: 24

LPAR Data, Collected in Partition VMLNX1

Processor type and model : 2094-606 Nr. of configured partitions: 7 Nr. of physical processors : 9 Dispatch interval (msec) : dynamic

Partition Nr	Upi	d #Proc	Weight	Wait-C	Сар	%Load	CPU	%Busy	%0vhd	%Susp	%VMI d	%LogI d	Туре
CF01A	1 0	1 1	900	NO	NO		0	98.7	98.7				I CF
CF01B	2 0	2 1	100	NO	NO		0	. 7	. 7				I CF
CPs removed for readability													
VMLNX1	6 1	3 2	500	NO	NO		0	47.2	47.2	1.9	46.7	47.6	I FL
			500		NO		1	47.0	47.0	1.9	46.6	47.5	I FL
VMLNX3	7 1	4 2	500	NO	NO		0	1.4	1.4				I FL
			500		NO		1	2.0	2.0				I FL
General LPAR mgmt overhead . 3													
Overall physical load 51.5													

If VMLNX3 ever gets hungry, there is going to be a CPU constraint.

_	
_	 _
_	_
_	
_	
_	 -

What We Recommended The Customer Change

More central storage

- Calculated new amount based on:
 - User pages resident on DASD ("before" FCX113 UPAGE report)
 - Understanding that memory comes in 32 GB increments

More and faster paging devices

- One would think if we were adding storage to reduce paging, we wouldn't have to tinker with the paging configuration too
- However, we knew the customer wanted to grow his logged-on virtual
- Also, customer reconfigured his [non-IBM] DASD to improve I/O response time

1 additional processor

- Seems sufficient based on old FCX225 SYSSUMLG report

-		
-	_	
_		
_		

Customer's New Configuration

	Old	New	Change
System model	2094-606	2094-705	-
Processors	2	3	+50%
SYSGEN storage	19968 (19.5 GB)	52736 (52 GB)	+267%
XSTORE	4096 (4 GB)	4096 (4 GB)	0
Page slots	24641k (94 GB)	72121k (275 GB)	+293%
Paging devices	25	40	+160%
Logged-on virtual	68 GB	87 GB	+28%



Measurement After The Change

- 60% reduction in user pages on DASD (FCX113 UPAGE)
- No queuing for paging devices (FCX109 DEVICE CPOWNED)
- No more user page waits (FCX114 USTAT)
- No processor constraint during the peak hour (FCX225 SYSSUMLG)

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AFTER: Customer's FCX109 DEVICE CPOWNED

FCX109 Run 2008/06/16 13: 5		EVICE CPOWN bad and Perfo		Owned Disks			
From 2008/06/13 11:35:25	E.			owned bi sks			SYSTEM
To 2008/06/13 15: 35: 25							CPU 20
For 14400 Secs 04:00:00		This is a per	formance rep	ort for SYSTE	M XY7"		z/VM
Page / SPOOL Allocation Su	ummary						
PAGE slots available	72121k	SPOOL slot	s available	600840			
PAGE slot utilization	4%	SPOOL slot	utilization	54%			
T-Disk cylinders avail.		DUMP slots	avai l abl e	0			
T-Disk space utilization	%	DUMP slot	utilization	%			
·							
< Device Descr>		<	Rate/s	>	User	Serv MLO	AD Block %Used
Volume Area	Area Used	<page></page>	<spool></spool>	SSCH	Inter <mark>Queu</mark>	<mark>e</mark> Time Res	sp Page for
Addr Devtyp Serial Type	Extent %	P-Rds P-Wrt	S-Rds S-Wrt	Total +RSCH	feres Lngt	<mark>h</mark> /Page Tir	me Size Alloc
790D 3390 520SPL SP00L	1-3338 54	. 0 . 0	1.5 1.6	3.0 3.1	1	0 2.1 2	. 1 100
7957 3390 520PAG PAGE	803060 4	.1.0		.1.1	1	0 3.0 3	.0
7958 3390 520PG3 PAGE	803060 4	.1.0		. 1 . 0	1	0 3.2 3	. 2
7959 3390 520PG5 PAGE	803060 4	.1.0		.1.0	1	0 4.2 4	. 2
795A 3390 520PG7 PAGE	803060 4	.1.0		.1.0	1	0 3.1 3	.1
795B 3390 520PG9 PAGE	803060 4	.1.0		.1.0	1	0 3.4 3	. 4
795C 3390 520PGB PAGE	803060 4	.1.0		.1.0	1	0 3.5 3	.5
795D 3390 520PGD PAGE	803060 4	.1.0		.1.0	1	0 4.1 4	.1
795E 3390 520PGF PAGE	803060 4	. 1 . 0		. 1 . 0	1	0 4.3 4	. 3
795F 3390 52PG11 PAGE	803060 4	. 1 . 0		. 1 . 0	1	0 3.7 3	. 7
7960 3390 52PG13 PAGE	803060 4	. 1 . 0		.1.1	1	0 3.3 3	. 3
69		89 WS		8/22/2009		© 200	7 IBM Corporation

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AFTER: Customer's FCX225 SYSSUMLG

FCX225 Run 2008/06/16 13: 51: 28	SYSSUMLG	
	System Performance Summary by Time	
From 2008/06/13 11:35:25		SYSTEMI D
To 2008/06/13 15: 35: 25		CPU 2094-7
For 14400 Secs 04:00:00	"This is a performance report for SYSTEM XYZ"	z/VM V.5

	<	CPU		>	<vec></vec>	<user< th=""><th>rs></th><th>< /</th><th>0></th><th><stg></stg></th><th><-Pagi</th><th>ng></th><th><spl></spl></th><th>< UP-</th><th>+MP Tra</th><th>nsacti</th><th>ons</th></user<>	rs>	< /	0>	<stg></stg>	<-Pagi	ng>	<spl></spl>	< UP-	+MP Tra	nsacti	ons
<rati o=""></rati>					SSCH	DASD	Users	<-Rate/s>			<-Response Time-> <-Transaction						
Interval	Pct		Cap-	0n-	Pct	Log-		+RSCH	Resp	in	PGI N+	Read+	Pages	Non-	Qui ck		Non-
End Time	Busy	T/V	ture	line	Busy	ged A	Activ	/s	msec	Elist	PGOUT	Write	/s	Triv Triv	Di sp	Triv	Tri v
>>Mean>>	14.1	1.06	. 9817	3.0		37	31	113.6	1.2	. 0	5.7	4.2	3.0	. 781 2. 788	. 959	3.84	1.80
11: 45: 25	14.0	1.05	. 9836	3.0		37	31	112.0	1.2	. 0	1.8	2.3	. 0	. 837 2. 980	. 963	3. 78	1.71
11: 55: 25	14.5	1.05	. 9838	3.0		37	31	104.5	1.3	. 0	1.2	3.7	. 0	. 821 2. 878	1.018	3.88	1.71
12: 05: 25	16. 1	1.05	. 9853	3.0		37	31	114.9	1.5	. 0	2.6	12.9	. 0	. 782 2. 807	1.040	3. 91	1.76
12: 15: 25	14.2	1.05	. 9835	3.0		37	31	106.7	1.3	. 0	1.8	4.3	. 0	. 789 2. 867	. 975	3.71	1.80
12: 25: 25	13.2	1.06	. 9822	3.0		37	31	112. 1	1.2	. 0	1.9	4.9	. 0	. 852 3. 054	. 953	3.71	1.64
12: 35: 25	13.6	1.05	. 9830	3.0		37	31	105.7	1.3	. 0	3.1	7.8	. 0	. 817 2. 687	. 887	3.99	1.77
12: 45: 25	14.1	1.05	. 9832	3.0		37	31	122.0	1.6	. 0	2.0	14.8	. 0	. 776 2. 996	. 967	3. 79	1.71
12: 55: 25	13.3	1.05	. 9825	3.0		37	30	102.6	1.3	. 0	1.1	3.3	. 0	. 783 2. 895	. 990	3.76	1.76
13: 05: 25	13.7	1.05	. 9828	3.0		37	31	114.9	1.2	. 0	1.1	5.2	. 0	. 787 2. 536	. 923	4.10	1.81
13: 15: 25	14.0	1.05	. 9832	3.0		37	30	104.5	1.2	. 0	. 6	1.1	. 0	. 783 2. 779	. 943	3.83	1.79
13: 25: 25	13.7	1.06	. 9826	3.0		37	30	114.1	1.2	. 0	. 9	4.0	. 0	. 756 2. 832	. 997	3.88	1. 78
13: 35: 25	13.6	1.06	. 9786	3.0		37	30	103.0	1.3	. 0	1.0	3.5	. 0	. 771 2. 676	. 990	3.77	1.88
13: 45: 25	13.7	1.05	. 9829	3.0		37	30	116. 1	1.1	. 0	. 9	2.0	. 0	. 793 2. 975	. 929	3.79	1.71
13: 55: 25	14.6	1.05	. 9838	3.0		37	30	102.2	1.2	. 0	. 8	4.2	. 0	. 749 2. 733	. 964	3.63	1.85
14: 05: 25	13.8	1.06	. 9830	3.0		37	31	115. 2	1. 1	. 0	. 6	2.7	. 0	. 766 2. 921	. 887	3.87	1. 78



AFTER: Customer's FCX126 LPAR

FCX126 Run 2008/06/16 13:51:28	LPAR
	Logical Partition Activity
From 2008/06/13 11:35:25	
To 2008/06/13 15: 35: 25	
For 14400 Secs 04:00:00	"This is a performance report for SYSTEM XYZ"

LPAR Data, Collected in Partition VMLNX1

Processor type and model	:	2094-705						
Nr. of configured partitions: 7								
Nr. of physical processors	:	9						
Dispatch interval (msec)	:	dynami c						

Partition Nr.	Upid #Pro	c Weight	Wait-C	Сар	%Load	CPU	%Busy	%0∨hd	%Susp	%VMI d	%LogI d	Туре
CF01A 1	01	1 900	NO	NO		0	98.2	98.2				I CF
CF01B 2	02	1 100	NO	NO		0	1.1	1.1				I CF
… CPs deleted for readability …												
VMLNX1 6	13	3 500	NO	NO		0	14.3	14.3	. 4	14.0	14.1	I FL
		500		NO		1	14.3	14.3	. 4	14.1	14.1	I FL
		500		NO		2	14.3	14.3	. 4	14.1	14.2	IFL
VMLNX3 7	14	3 500	NO	NO		0	2.9	2.9				I FL
		500		NO		1	2.9	2.9				I FL
		500		NO		2	2.9	2.9				I FL
General LPAR m	gmt overhea	ad			. 4							

41.0

Overall physical load



Summary

- System was running "as it was designed to run" before the changes
- Workloads ran as expected after the changes

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Case Study: PAV and MDC

8/22/2009

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Customer Called IBM

I have disk I/O problems on z/VM 5.2

When I turn on MDC, my system slows down

•OK, nobody panic or speculate

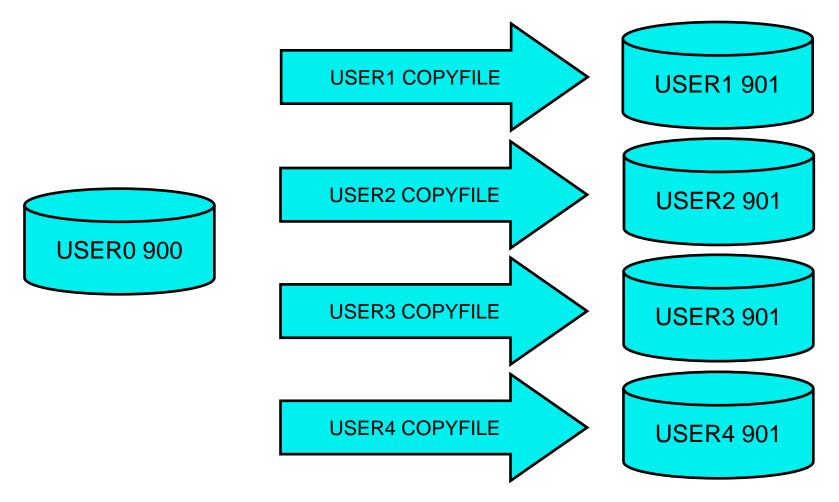
Send us some raw monitor data...

- ► For MDC off, when things are good
- ► For MDC on, when things are not so good
- Customer sent two very descriptive sets of data

So good, in fact, that we easily replicated the customer's problem on GDLSPRF3



Customer's Workload



All five minidisks are on the same RDEV.

8/22/2009

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Without MDC

z/VM 5.2, MDC OFF, without the fix, excerpt from typical FCX108 (DEVICE) report

< Device Descr>	Mdisk Pa-	<-Rate/s->	< Time	(msec)> Req.	
Addr Type Label/ID	Links ths	I/O Avoid	Pend Disc Con	n Serv Resp CUWt Qued	
E700 3390 LDB307	15 4	564 .0	.3 .2 1.	2 1.7 1.7 .0 .0	
E7FC ->E700 LDB307	15 4	543.0	.3 .2 1.	2 1.7 1.7 .0 .0	
E7FD ->E700 LDB307	15 4	541 .0	.4 .2 1.	2 1.8 1.8 .0 .0	
E7FE ->E700 LDB307	15 4	539.0	.4 .2 1.	2 1.8 1.8 .0 .0	

Things to notice:

- E700 with three PAV aliases
- Four users doing I/O to the volume
- Aggregate volume I/O rate is (564+543+541+539) = 2187 IOs/sec

From this report, how do we know...

- PAV is correctly configured for the volume?
- PAV is functioning correctly?
- MDC is turned off for this volume?



But With MDC...

z/VM 5.2, MDC ON, without the fix

< Device I	Descr>	Mdisk	Pa-	<-Rat	ce/s->	<	:	ſime	(msec))	>	Req.
Addr Type	Label/ID	Links	ths	I/O	Avoid	Pend	Disc	Conn	Serv	Resp	CUWt	Qued
E700 3390	LDB307	15	4	564	90.6	.3	.1	1.3	1.7	1.9	.0	2.9
E7FC ->E700	LDB307	15	4	.5	.0	.3	.3	1.5	2.1	2.3	.0	.0
E7FD ->E700	LDB307	15	4	.3	.0	.3	.3	1.4	2.0	2.2	.0	.0
E7FE ->E700	LDB307	15	4	.3	.0	.3	.6	1.3	2.2	2.4	.0	.0

Things to ask ourselves:

- Is MDC really on for this volume?
- What is the MDC-on volume I/O rate? (564+90.6+.5+.3+.3) = 655.7 IOs/sec
- Is it correct that the volume I/O rate should go down with MDC on?
- Is the CP I/O subsystem functioning correctly? Why or why not?



Things That Look Suspicious

- The volume I/O rate should not go down substantially when MDC comes online.
 - If some other limit is holding the applications back, the volume I/O rate should stay about the same
 - If nothing else is holding the applications back, the volume I/O rate should increase
- The CP I/O subsystem appears not to be functioning correctly
 - ► I/Os are happening almost exclusively on the base RDEV
 - ► The aliases appear to be doing almost no work
 - ► There is queueing at the base RDEV



A Visit to Development

I went to see Bill Stephens (virtual I/O and MDC expert)

He felt MDC's I/Os should be PAV-able

But investigation revealed...

79

- There are bits CP sets for its own I/Os to tell the real I/O layer whether to try to PAV the I/O...
- ► but MDC was forgetting to set these bits...
- ► (in fact, nowhere did CP ever set those bits!)...
- ► thus all I/Os originating in MDC were being forced to the base...
- ► thus MDC was failing to exploit the volume's PAV capability



MDC ON, With The Fix

z/VM 5.2, MDC ON, with the fix

< Device Descr>	Mdisk Pa-	<-Rate/s->	< Time	(msec)> Req.
Addr Type Label/ID	Links ths	I/O Avoid	Pend Disc Conn	Serv Resp CUWt Qued
E700 3390 LDB307	15 4	442 402.3	.3 .4 1.4	2.1 2.1 .0 .0
E7FC ->E700 LDB307	15 4	421 .0	.4 .3 1.5	2.2 2.2 .0 .0
E7FD ->E700 LDB307	15 4	415 .0	.4 .3 1.5	2.2 2.2 .0 .0
E7FE ->E700 LDB307	15 4	410 .0	.4 .3 1.5	2.2 2.2 .0 .0

Things to notice:

- MDC is functioning (there are avoided I/Os)
- Aggregate I/O rate is (442+402+421+415+410) = 2090 IOs/sec
- About one-fifth of the I/Os are being avoided... makes sense
- Connect time is up compared to MDC off (1.2 to 1.5) -- IOs are bigger
- I/Os are spreading across base and aliases
- No queueing at the base device



It Turns Out...

- Our System Test group saw this behavior too, during z/VM 5.2 test
- They thought it was expected that the I/O rate would go down when MDC was ON, and we can't entirely blame them
- The only clue anything is wrong is that there is a queue at the base RDEV -- I doubt System Test would see that
- All of our PAV measurements were done with MDC OFF, of course



Status

=APAR VM64199, UM32047 (z/VM 5.2), UM32048 (z/VM 5.3)

Is on the GA RSU for z/VM 5.3

In the base of z/VM 5.4

8/22/2009

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Case Study: Paging Difficulties

8/22/2009



Customer Calls In

- My system isn't running fast, but it isn't paging either
- •My application formats lots of VDISKs... aren't they in memory? Shouldn't this be fast?
- I have raw monitor data... will you take a look?
- Customer sent raw monitor file 20070501 MD111606
- •He says his workload uses disk volumes 1240-59 and 16C0-E3
- We took a look-see

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Basic System Summary

FCX225 Run 2007/05/02 12: 56: 34

SYSSUMLG

System Performance Summary by Time

From 2007/05/0111:16:08To2007/05/0112:37:10For4861Secs01:21:01

Result of 20070501 Run

	<	CPU		>	<vec></vec>	<use< th=""><th>ers></th><th><1/</th><th>0></th><th><stg></stg></th><th><-Pagi</th><th>ng></th></use<>	ers>	<1/	0>	<stg></stg>	<-Pagi	ng>
		<rat< td=""><td>i o></td><td></td><td></td><td></td><td></td><td>SSCH</td><td>DASD</td><td>Users</td><td><-Rate</td><td>e/s></td></rat<>	i o>					SSCH	DASD	Users	<-Rate	e/s>
Interval	Pct		Cap-	0n-	Pct	Log-		+RSCH	Resp	in	PGI N+	Read+
End Time	Busy	T/V	ture	line	Busy	ged	Activ	/s	msec	Elist	PGOUT	Write
>>Mean>>	10. 3	106.3	. 7577	27.0		280	263	122.7	11. 1	. 0	5418	1445
11: 23: 41	9.9	180. 7	. 8232	27.0		280	263	25.8	. 8	. 0	2645	. 0
11: 24: 40	10. 3	193 . 5	. 8051	27.0		280	263	23.8	. 7	. 0	2707	. 0
11: 25: 39	10. 5	196.8	. 8218	27.0		280	262	23.6	. 8	. 0	2825	. 0
11: 27: 10	9.7	159.5	. 8232	27.0		280	262	29.9	. 7	. 0	3714	. 0
11: 28: 09	9.8	108. 2	. 8015	27.0		280	266	48.4	. 8	. 0	8942	. 1
11: 29: 40	9.8	119. 2	. 8134	27.0		280	264	33.2	. 9	. 0	8602	2.8
11: 36: 10	10. 3	119.6	. 8048	27.0		280	263	45. 7	. 6	. 0	9327	. 0
11: 37: 40	10.5	136.8	. 8028	27.0		280	262	30.3	. 6	. 0	9213	. 0
11: 39: 10	10. 8	144.2	. 8158	27.0		280	264	30.7	. 7	. 0	9189	. 0
11: 40: 40	10.5	135.6	. 8093	27.0		280	264	32.5	. 7	. 0	10083	. 0
11: <u>41: 39</u>	10.7	166_5	8124	27_0		280	262	25_2		0	8942	. 0
	ok at	tho	se T	Γ/\/	rati	lan	Wh	at is	CF	op do	ind	. 0
	n al			i / V	iui		VVI			uu	<u>inig</u> :	



Think About the Application

- Customer says he is formatting VDISKs
- VDISKs are address spaces
- •We page them when storage gets tight
- •We do seem to be spending a lot of time in CP
- Let's see if DEVICE CPOWNED shows us anything

0-3338

2



DEVICE CPOWNED

FCX109 Run 2007/05/02 12: 56: 34 DEVICE CPOWNED Load and Performance of CP Owned Disks											
From 2007/05/01 11:16:08										-	20070501
To 2007/05/01 12: 37: 10										(CPU 2094
For 4861 Secs 01: 21: 01	Resul t	of 20070501	Run							Z	z/VM V
Page / SPOOL Allocation Summary											
PAGE slots available 34745k	SPC	OL slots av	ai I abl e	36	56598						
PAGE slot utilization 3%	SPC	OL slot uti	lization		9%						
T-Disk cylinders avail.	DUN	P slots ava	ilable		0						
T-Disk space utilization%	DUM	P slot util	ization		%						
·											
< Device Descr>	<		Rate/s -		>	User		Serv	MLOAD	BI ock	%Used
Volume Area Area	Used <p< td=""><td>age> <</td><td>Spool></td><td></td><td>SSCH</td><td>Inter</td><td>Queue</td><td>Time</td><td>Resp</td><td>Page</td><td>for</td></p<>	age> <	Spool>		SSCH	Inter	Queue	Time	Resp	Page	for
Addr Devtyp Serial Type Extent	% P-Ro	s P-Wrt S-R	ds S-Wrt	Total	+RSCH	feres	Lngth	/Page	Time	Si ze	Alloc
1240 3390 XXPG20 PAGE 0-3338	3 1.	2 17.6 .		18. 8	1.4	1	0	3.8	3.8	14	44
1241 3390 XXPG21 PAGE 0-3338	3 1.	3 16.8 .		18. 1	1.3	1	0	7.8	7.8	14	42
1242 3390 XXPG22 PAGE 0-3338	3 1.	3 17.4 .		18.6	1.3	1	. 57	6.7	9.0	14	43
1243 3390 XXPG23 PAGE 0-3338	2 1.	3 16.2 .		17.5	1.3	1	1.08	5.2	11.0	14	40
1244 3390 XXPG24 PAGE 0-3338	2 1.	3 16.4 .		17.7	1.3	1	1. 16	5.0	11.5	14	41
1245 3390 XXPG25 PAGE 0-3338	2 1.	2 15.9 .		17.1	1.3	1	. 57	5. 6	8.6	14	40

From 11:16 to 12:37 the paging devices have queues on average? Let's look at some INTERIM reports and see what we see...

1.3 15.7

1.2

1

14

39

0 12.5 12.5

....

1246 3390 XXPG26 PAGE



INTERIM DEVICE, 11:47

1FCX108 Run 2007/05/02 12: 56: 29

INTERIM DEVICE General I/O Device Load and Performance

From2007/05/0111: 45: 39To2007/05/0111: 47: 37For118Secs00: 01: 58

Result of 20070501 Run

< Device	Descr>	Mdi sk	Pa-	<-Rat	te/s->	<	7	Γ im e	(msec)		>	Req.	<perc< td=""><td>:ent></td></perc<>	:ent>
Addr Type	Label /I D	Li nks	ths	I/0	Avoi d	Pend	Di sc	Conn	Serv	Resp	CUWt	Qued	Busy	READ
1240 3390	XXPG20 CP	0	2	1.3	. 0	47.3	. 9	5.4	53.6	53. <mark>6</mark>	. 0	. 0	14	0
16DE 3390	XXPGOE CP	0	2	1.3	. 0	48.9	. 1	6.5	55.5	55.5	. 0	. 0	12	100
16E0 3390	XXPGOF CP	0	2	1.3	. 0	53.8	. 6	7.0	61.4	61.4	. 0	. 0	12	0
16D9 3390	XXPGOD CP	0	2	1.3	. 0	53.3	. 9	6. 1	60.3	60.3	. 0	. 0	14	100
16DF 3390	XXPG09 CP	0	2	1.3	. 0	49.9	. 0	7.1	57.0	57.0	. 0	. 0	11	100
16DC 3390	XXPG07 CP	0	2	1.2	. 0	50.7	. 0	6.5	57.2	57.2	. 0	. 0	12	100
1247 3390	XXPG27 CP	0	2	1.2	. 0	52.2	. 7	6.4	5 9 . 3	75.0	. 0	. 0	15	0
16DB 3390	XXPG06 CP	0	2	1. 2	. 0	51.6	. 0	7.0	58. 6	58. <mark>6</mark>	. 0	. 0	12	0
16DD 3390	XXPG08 CP	0	2	1.2	. 0	54.6	. 4	7.2	62. 2	62. 2	. 0	. 0	13	0
16D8 3390	XXPGOC CP	0	2	1. 2	. 0	54.7	. 0	6.6	61.3	61.3	. 0	. 0	13	100
1241 3390	XXPG21 CP	0	2	1. 2	. 0	48.9	. 8	7.0	56.7	56.7	. 0	. 0	13	0
16D6 3390	XXPGOB CP	0	2	1.1	. 0	55.7	. 5	6.9	63.1	63.1	. 0	. 0	13	0
1242 3390	XXPG22 CP	0	2	1.1	. 0	45.5	. 0	7.3	52.8	52.8	. 0	. 0	12	0

Look at that pending time on the paging volumes! High pending time usually means channel contention...



Configuration

From FCX131 DEVCONF:

1240-1259	0008-0021	3390-3 (E)	67 69				2105-E8	0nl i ne
16C0-16E3	0050-0073	3390-3 (E)	67 69				2105-E8	0nl i ne

From FCX161 LCHANNEL:

67	ESCON	00 Yes	15	6	93	7
69	ESCON	00 Yes	16	10	73	27

Two ESCON chpids for all this paging DASD?!



Recommendation

Customer added four ESCON chpids

•(Why didn't he add FICON?

He was quiet for a while, and then...

IBM Systems and Technology Group



AB815 CPU 209

z/VM

He's Baa-aaack

FCX109 Run 2007/08/15 09: 58: 19

INTERIM DEVICE CPOWNED Load and Performance of CP Owned Disks

From	2007/08/14	07: 15: 03
То	2007/08/14	07: 20: 02
For	299 Secs	00: 04: 59

Result of AB815 Run

· · · ·											
< Device Descr>		< R	ate/s -		>	User		Serv	MLOAD	BI ock	%Used
Volume Area	Area Used	<page> <s< td=""><td>pool></td><td></td><td>SSCH</td><td>Inter</td><td>Queue</td><td>Time</td><td>Resp</td><td>Page</td><td>for</td></s<></page>	pool>		SSCH	Inter	Queue	Time	Resp	Page	for
Addr Devtyp Serial Type	Extent %	P-Rds P-Wrt S-Rd	s S-Wrt	Total	+RSCH	feres	Lngth	/Page	Time	Si ze	Alloc
16D5 3390 XXPGOA PAGE	0-3338 88	21.7 19.1		40.8	15.5	1	33.00	1.0	2.9	4	49
16D6 3390 XXPGOB PAGE	0-3338 88	20.5 17.2		37.7	15.1	1	19.00	2. 2	42.5	4	44
16D8 3390 XXPGOC PAGE	0-3338 88	22.7 18.1		40.7	15.8	1	22.00	1. 2	28.7	4	45
16D9 3390 XXPGOD PAGE	0-3338 87	21.1 18.5		39.6	15.2	1	29.00	. 8	25.0	4	48
16DB 3390 XXPG06 PAGE	0-3338 87	22.3 20.0		42.3	15.6	1	20.00	. 8	17.1	4	51
16DC 3390 XXPG07 PAGE	0-3338 86	21.9 17.7		39.6	15.7	1	10.00	. 9	10.4	3	45
16DD 3390 XXPG08 PAGE	0-3338 86	22.0 18.2		40.3	15.5	1	106. 0	. 8	5. 9	4	47
16DE 3390 XXPGOE PAGE	0-3338 86	21.4 19.6		41.0	15.0	1	0	. 6	. 6	4	48
16DF 3390 XXPG09 PAGE	0-3338 84	22.1 19.6		41.7	14.2	1	17.00	1.0	18.4	5	50
16E0 3390 XXPGOF PAGE	0-3338 83	20.4 17.6		38.1	12.4	1	63.00	2.2	139.3	5	44
5805 3390 CF5805 PAGE	810000 12	46.5 41.9		88.4	21.3	10	0	. 1	. 1	11	100
9F23 3390 XPG2 PAGE	0-3338 99	18.6 18.1		36.7	25. 9	1	23.00	. 7	16. 9	2	47
9F24 3390 XPG3 PAGE	0-3338 99	19.2 17.5		36.6	25.8	1	29.00	. 6	19. 2	2	46
9F25 3390 XPG4 PAGE	0-3338 99	18.6 17.4		36.0	26.9	1	0	. 6	. 6	1	46
9F2F 3390 XPG6 PAGE	0-3338 99	20.9 17.9		38.8	27.1	1	35.00	. 6	20.6	2	47
CO9E 3390 PCO9B PAGE	0-3338 100	22.4 19.2		41.6	30. 2	1	0	. 6	. 6	1	98
D007 3390 CFD007 PAGE	896800 17	46.1 40.7		86.8	19.9	1	30.00	. 1	. 1	11	99
D008 3390 CFD008 PAGE	896800 17	42.2 39.7		81.9	18.1	1	32.00	. 2	. 2	11	99
DOOD 3390 JSPG04 PAGE	896800 20	42.9 39.0		81.9	18.5	1	0	. 3	. 3	12	100

Removed 25 100% full 3990-3's from this excerpt.

8/22/2009



So What's The Problem

- •40 3390-3 paging volumes nearly full
- •4 3390-9 paging volumes have the free space
- We can do only one I/O at a time to those gigantic model 9's
- •Get rid of those mod 9's and add a lot of mod 3's



Case Study: HiperSockets

8/22/2009



Question from Customer

My system seems to have:

- Long ping times from my z/VM partition to my z/OS partition
- Long transaction times from my z/VM partition to my z/OS partition
- Seems related to my use of HiperSockets to connect z/VM to z/OS
 - When I use a real OSA to connect the partitions, I don't have these problems
 - When I drive the z/OS server from external AIX boxes, I don't have the long transaction times

Customer sent MONWRITE data



CEC and LPAR Configuration, from MONWRITE Data

- 2094 with:
 - 6 CPs
 - 4 ICFs
 - 12 IFLs
 - 4 zIIPs
- Several z/VM partitions, all shared IFL 12-ways, but only one of these partitions is active
- Several z/OS partitions, all shared, with varying logical PU configurations, that use the CPs and zIIPs
- Two coupling partitions that use the ICFs, dedicated



Workload Configuration, from Customer

- Linux on z/VM is the origin
- Over to z/OS database server via real HiperSocket
- Back to Linux on z/VM via real HiperSocket
- Transaction ends

_			
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		_	
	_	_	
_	-	-	
	_	_	
_			

Long Response Time? Let's Hunt z/VM Constraints

- FCX126 LPAR and FCX225 SYSSUMLG no IFL constraints found
- FCX225 SYSSUMLG and FCX109 INTERIM DEVICE CPOWNED no paging found – everything fits in central
- FCX108 DEVICE DASD revealed the active user volumes, and FCX168 DEVLOG showed good service time and small to no queues
- FCX215 INTERIM FCHANNEL no FICON adapter CPU problems; FCX108 DEVICE DASD – no pending time concerns
- FCX112 USER revealed the big CPU users, and FCX162 USERLOG showed very low T/V and no CPU peaks
- FCX231 INTERIM HIPSOCK showed <10 msgs/sec and 600 data units/message – seems small
- Couldn't find a z/VM constraint



FCX126 LPAR: General View of CPU Busy (z/VM)

Parti ti on	Nr.	Upi d	#Proc	Weight	Wait-C	Сар	%Load	CPU	%Busy	%0vhd	%Susp	%VMI d	%LogId T	уре
XXX1	11	14	12	200	NO	NO		0	20.5	. 1	. 3	20.3	20.4 I	FL
				200		NO		1	21.9	. 2	. 3	21.7	21.7 I	FL
				200		NO		2	21.5	. 2	. 3	21.3	21.3 I	FL
				200		NO		3	22.0	. 2	. 3	21.7	21.8 I	FL
				200		NO		4	22. 1	. 2	. 3	21.9	21.9 I	FL
				200		NO		5	22.2	. 2	. 3	22.0	22.0 I	FL
				200		NO		6	22.0	. 2	. 3	21.8	21.8 I	FL
				200		NO		7	22. 1	. 2	. 3	21.8	21.9 I	FL
				200		NO		8	21.9	. 1	. 3	21.7	21.8 I	FL
				200		NO		9	22.2	. 2	. 3	22.0	22.0 I	FL
				200		NO		10	22. 1	. 2	. 3	21.8	21.9 I	FL
				200		NO		11	21.9	. 2	. 3	21.6	21.7 I	FL

The other partitions using IFLs are not running. We don't see a problem here.



FCX126 LPAR: CPU Busy, z/OS Partitions

Partition	Nr.	Upi d	#Proc	Weight	Wait-C	Сар	%Load	CPU	%Busy	%0vhd	%Susp	%VMI d	%LogI d	Туре
XXXA	8	03	8	190	NO	NO	5. 9	0	28.2	. 2				СР
				190		NO		1	28.2	. 2				СР
				190		NO		2	28. 1	. 2				СР
				190		NO		3	28.0	. 2				СР
				190		NO		6	10.6	. 1				ZIIP
				190		NO		7	10.5	. 0				ZIIP
				190		NO		22	10.5	. 0				ZIIP
				190		NO		23	10.6	. 1				ZIIP
XXXB	9	04	8	190	NO	NO	5.8	0	26.5	. 1				СР
				190		NO		1	26.5	. 2				СР
				190		NO		2	26.4	. 1				СР
				190		NO		3	26.4	. 1				СР
				190		NO		6	11.5	. 1				ZIIP
				190		NO		7	11.5	. 0				ZIIP
				190		NO		22	11.5	. 0				ZIIP
				190		NO		23	11.5	. 0				ZIIP
XXXO	10	05	2	20	NO	NO		0	. 0	. 0				СР
				20		NO		1	. 0	. 0				ZIIP
YYYA	13	11	9	600	NO	NO	23.5	0	70.5	. 1				СР
				600		NO		1	70.5	. 1				СР
				600		NO		2	70.4	. 1				СР
				600		NO		3	70.2	. 1				СР
				600		NO		4	64.9	. 0				ZIIP
				600		NO		5	64.9	. 0				ZIIP
				600		NO		6	69.3	. 1				СР
				600		NO		22	64.9	. 0				ZIIP
				600		NO		23	65.0	. 1				ZIIP

None of these partitions look real busy. But we will see

shortly that there is in fact a problem.



How Do HiperSockets work?

- Synchronous data transfer between partitions
- Firmware copies data from one PU's buffers to the other PU's buffers
- Firmware runs on the PUs that did the calls
- Works fine in CPU-rich environments
- Not so well in CPU-constrained environments

_	
_	
	 = 7 =

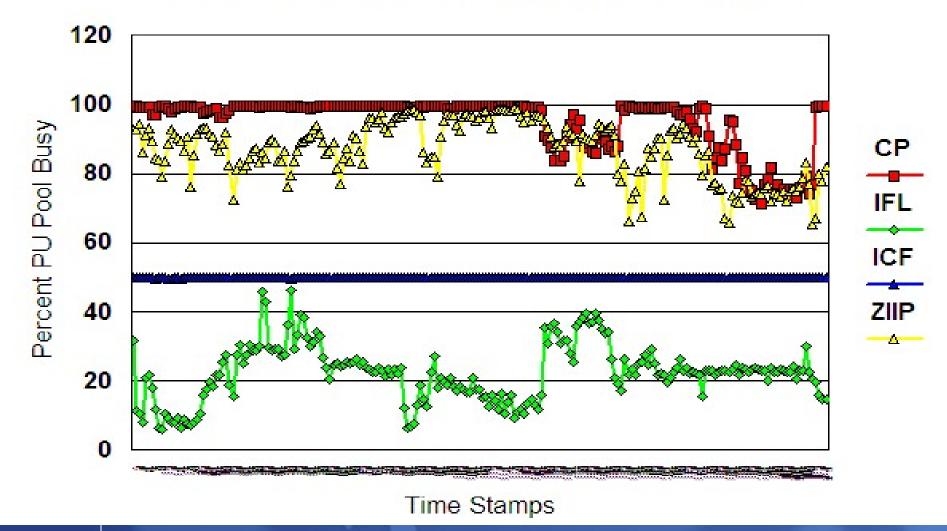
How Are The z/OS CPUs Doing?

- 6 real standard CPs
- 4 z/OS partitions have 14 logical standard CPs altogether
- Good thing MONWRITE data captures LPAR's view of the partitions' consumptions
- For each sample interval, add up those 14 logical standard CPs' utilizations to see how much of the 6 real CPs they're using altogether
 - Post-process the FCX126 INTERIM LPAR reports
 - Requires a little Rexx
- While we're at it, do this for every engine type



CEC View, PU Utilization by Type

D090408, PU Busy by PU Type



102

8/22/2009



Findings

- Real standard CPs are saturated
- Real HiperSockets require readily accessible cycles to work well
- This is a CPU-constrained environment
- Explains why real OSA was better
- Recommended either:
 - Adding more standard CPs, or
 - Tuning z/OS partitions to reduce CP resource they need
- Also recommended a z/OS expert look at RMF reports to see how else the CP constraint was impacting z/OS function

_	

Case Study: Grinch that Stole Performance



The Grinch That Stole Performance

From	Perfo	ormanc	e Too	olkit	USTAT	FCX1	L14 I	Repor	rt Ja	anuary	5:
<-SVM and->											
%CPU	%LDG	%PGW	%IOW	%SIM	%TIW	%CFW	%TI	%EL	%DM	%IOA	
0	0	0	19	2	10	0	3	0	51	8	

										ary 5:
	<-Rat	ce/s->	<	[rime ((msec))	>	Req.	<pct></pct>
Addr	I/O	Avoid	Pend	Disc	Conn	Serv	Resp	CUWt	Qued	Busy
1742	26.7	.0	1.3	18.4	4.7	24.5	69.0	.0	1.2	65.4

Went to check Toolkit CACHEXT FCX177 Report for control unit cache stats, but it didn't exist!

It is a good thing I keep historical data -- let's go back and see what's going on...



When Did We Last See Cache?

From	Perfo	ormance	e Tool	lkit 1	DEVICE	E FCX1	108 Re	eport	:	
	<-Rat	ce/s->	<	'	Time ((msec))	>	Req.	<pct></pct>
Addr	I/O	Avoid	Pend	Disc	Conn	Serv	Resp	CUWt	Qued	Busy
Dec8	41.0	.0	0.3	0.2	2.0	2.6	2.9	.0	.0	10.5
Jan5	26.7	.0	1.3	18.4	4.7	24.5	69.0	.0	1.2	65.4

From	Perform	mance	Toolkit	CACHE	EXT FO	CX177	7 Dec	. 8 th	Rep	ort:
< Rate/s> <percent></percent>										
Total	Total	Read	l Read	Write		<]	Hits		>
Cache	SCMBK	N-Sec	I Seq	FW	Read	Tot	RdHt	Wrt	DFW	CFW
53.0	41.0	52.3	8 0	0.6	99	99	99	96	96	• •



Down for the 3-Count

q dasd details 1742 1742 CUTYPE = 3990-EC, DEVTYPE = 3390-06, VOLSER=USE001 CACHE DETAILS: CACHE NVS CFW DFW PINNED CONCOPY -SUBSYSTEM F Y Y Y Ν _ -DEVICE Υ Υ Ν Ν — DEVICE DETAILS: CCA = 02, DDC = 02DUPLEX DETAILS: SIMPLEX

Pinned data! Yikes! I had never seen that before!



Performance Toolkit Device Details

FCX110 CPU 20	03 GDL	VM7 Int	terval	. IN	ITIAL	13:08:47	Remote	Data
		1 1 1 1 4 0	/ 0170					
Detailed Analysi	s Ior De	vice 1/42	(SYS	2.T.F.M)			
Device type : 3	390-2	Function	pend.	:	.8ms	Device bu	isy :	27%
VOLSER : U	SE001	Disconne	cted	:	20.3ms	I/O conte	ention:	0%
Nr. of LINKs:	404	Connected	f	:	5.4ms	Reserved	:	0%
Last SEEK :	1726	Service (time	:	26.5ms	SENSE SSC	CH :	•••
SSCH rate/s :	10.5	Response	time	:	26.5ms	Recovery	SSCH :	•••
Avoided/s :	• • • •	CU queue	time	:	.Oms	Throttle	del/s:	• • •
Status: SHARABLE								
Path(s) to devic	e 1742:	0A 2	2A	4A				
Channel path sta	tus :	ON (ON	ON				
Device	Overall	CU-Cache	Perfo	orma	nce	Split	2	
DIR ADDR VOLSER	IO/S %RI	EAD %RDHI	Г %WRH	IIT	ICL/S BY	P/S IO/S	%READ %I	RDHIT
08 1742 USE001	.0	0 (C	0	.0	.0 'NORN	MAL' I/O	only



Per	form	ance T	oolk	it Dev	vice	Details
MDISK	Extent	Userid	Addr	Status	LINK	MDIO/s
101 -	- 200	EDLSFS	0310	WR	1	.0
201 -	- 500	EDLSFS	0300	WR	1	.0
501 -	- 600	EDLSFS	0420	WR	1	.0
601 -	- 1200	EDLSFS	0486	WR	1	.0
1206 -	- 1210	RAID	0199	owner		
		BRIANKT	0199	RR	5	.0
1226 -	- 1525	DATABASE	0465	owner		
		K007641	03A0	RR	3	.0
1526 -	- 1625	DATABASE	0269	owner		
		BASILEMM	0124	RR	25	.0
1626 -	- 1725	DATABASE	0475	owner		
		SUSANF7	0475	RR	1	.0
1726 -	- 2225	DATABASE	0233	owner	366	10.5



Solution

Use Q PINNED CP command to check for what data is pinned.

- Discussion with Storage Management team.
- Moved data off string until corrected.

Pinned data is <u>very</u> rare, but when it happens it is serious.