



IBM Systems & Technology Group

z/VM System Limits

Session 9160

IBM z/VM Performance Evaluation
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Agenda

- **Describe various limits**
 - Architected
 - Consumption
 - Latent
- **Show how to keep tabs on consumables**
- **Discuss limits that may be hit first**

Limits

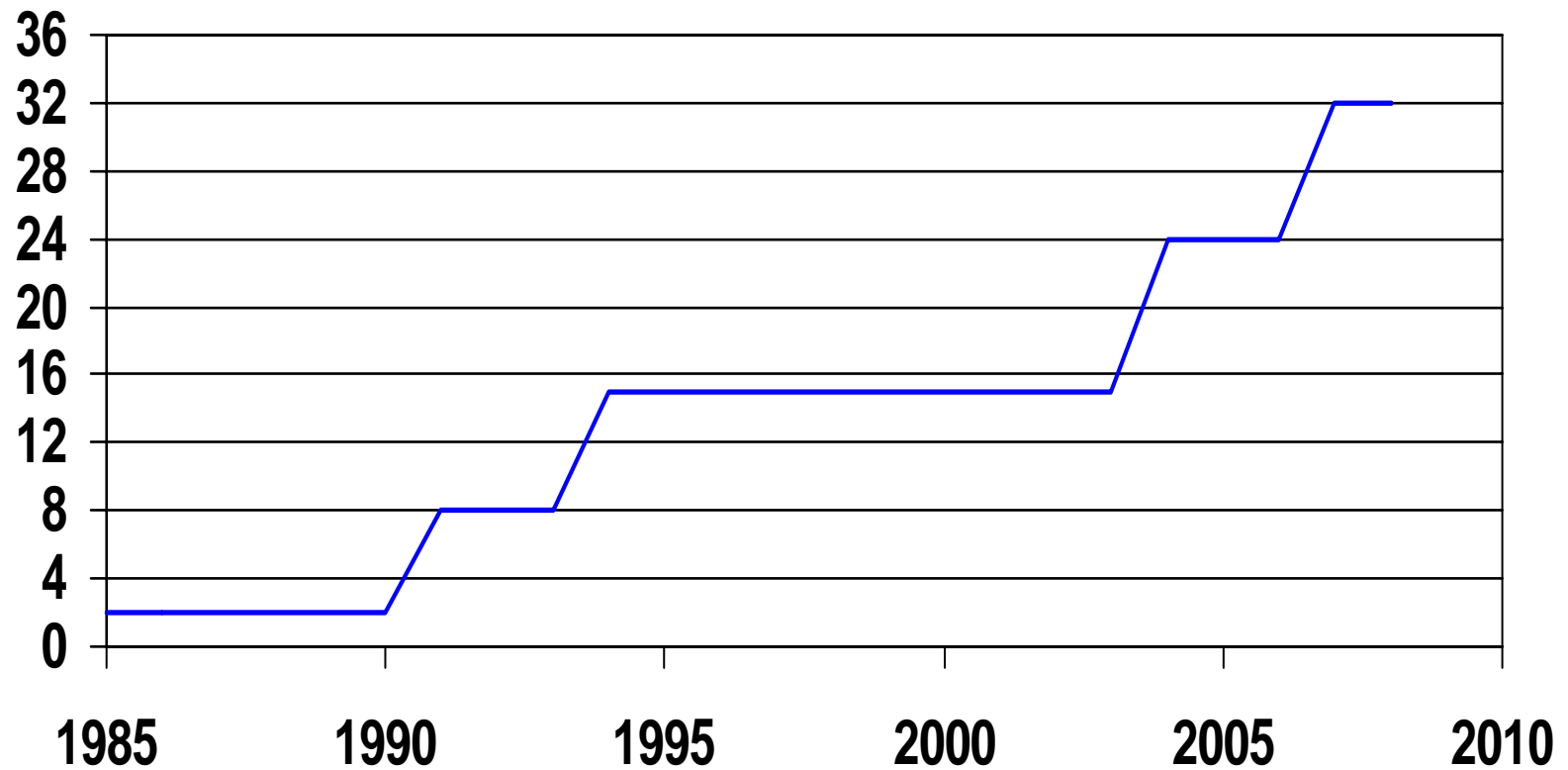
- **Processors**
- **Memory**
- **I/O**
- **Others**
- **Latent limits**

Processors

- **Processors (architected): 64**
 - Includes all engine types (CP, zAAP, zIIP, IFL...)
- **Processors (hardware): 64 (z10 EC); 54 (z9 EC)**
- **Logical processors (unsupported): 64 (z10 EC); 54 (z9 EC)**
- **Logical processors (support statement): 32**
- **Master processor (architected): 1**
 - 100%-utilized master is the issue
 - z/VM will elect a new master if master fails
- **Virtual processors in single virtual machine (architected): 64**
 - But $N_{\text{Virtual}} > N_{\text{Logical}}$ is not usually practical
- **Number of partitions: 60 (z9 EC & z10 EC)**

Processor Scaling

Number of Supported Processors



Processors: FCX100 CPU

FCX100 Run 2007/09/06 14:00:28

CPU

General CPU Load and User Transactions

From 2007/09/04 09:07:00

To 2007/09/04 10:00:00

For 3180 Secs 00:53:00

CPU 2094-700

z/VM V.5.3.0 SLU 0701

CPU Load										Vector Facility			Status or	
PROC	TYPE	%CPU	%CP	%EMU	%WT	%SYS	%SP	%SIC	%LOGLD	%VTOT	%VEMU	REST	ded.	User
P00	I FL	16	2	14	84	2	0	84	16	
P15	I FL	18	2	16	82	1	0	80	18	
P14	I FL	18	2	16	82	1	0	80	18	
P13	I FL	18	2	16	82	1	0	80	18	
P12	I FL	18	2	16	82	1	0	81	18	
P11	I FL	18	2	17	82	1	0	80	19	
... truncated ...														

1. $T/V \sim 18/16 = 1.13$ a little CP overhead here
2. Master does not seem unduly burdened

Processors: FCX114 USTAT

FCX114 Run 2007/09/06 14:00:28

USTAT

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Wait State Analysis by User

From 2007/09/04 09:07:00

To 2007/09/04 10:00:00

For 3180 Secs 00:53:00

CPU 2094-700

z/VM V.5.3.0 SLU 0701

User id	%ACT	%RUN	%CPU	%LDG	%PGW	%LOW	%SIM	<-SVM and->								<--%Time spent in-->					Nr of Users	
								%TIW	%CFW	%TI	%EL	%DM	%IOA	%PGA	%LIM	%OTH	Q0	Q1	Q2	Q3		E0-3
>System<	64	1	0	1	0	0	0	83	0	0	0	3	0	0	0	10	1	29	10	57	0	211
TCPIP	100	0	0	0	0	0	0	0	0	3	0	97	0	0	0	0	3	0	0	0	0	0
RSCSDNS1	100	0	0	0	0	0	0	0	0	0	0	100	0	0	0	0	0	0	0	0	0	0
SNMPD	100	0	0	0	0	0	0	0	0	2	0	98	0	0	0	0	2	0	0	0	0	0
SZVAS001	100	2	0	0	0	0	0	97	0	0	0	0	0	0	0	1	0	3	12	85	0	0

1. %CPU wait is very low – nobody is starved for engine
2. %TIW is “test idle wait” – we are waiting to see if queue drop happens

Memory

■ **Central storage**

- Supported central storage: 256 GB
- Unsupported central storage:
 - 512 GB minus your HSA (z9 EC)
 - 1 TB (z10 EC)
- The largest we ever managed was 440 GB

■ **Expanded storage (architected): 16TB**

- z/VM Limit: 128GB
- See <http://www.vm.ibm.com/perf/tips/storconf.html>

■ **Virtual machine size (hardware):**

- On z9 and z10, 1 TB (2^{40})

Memory

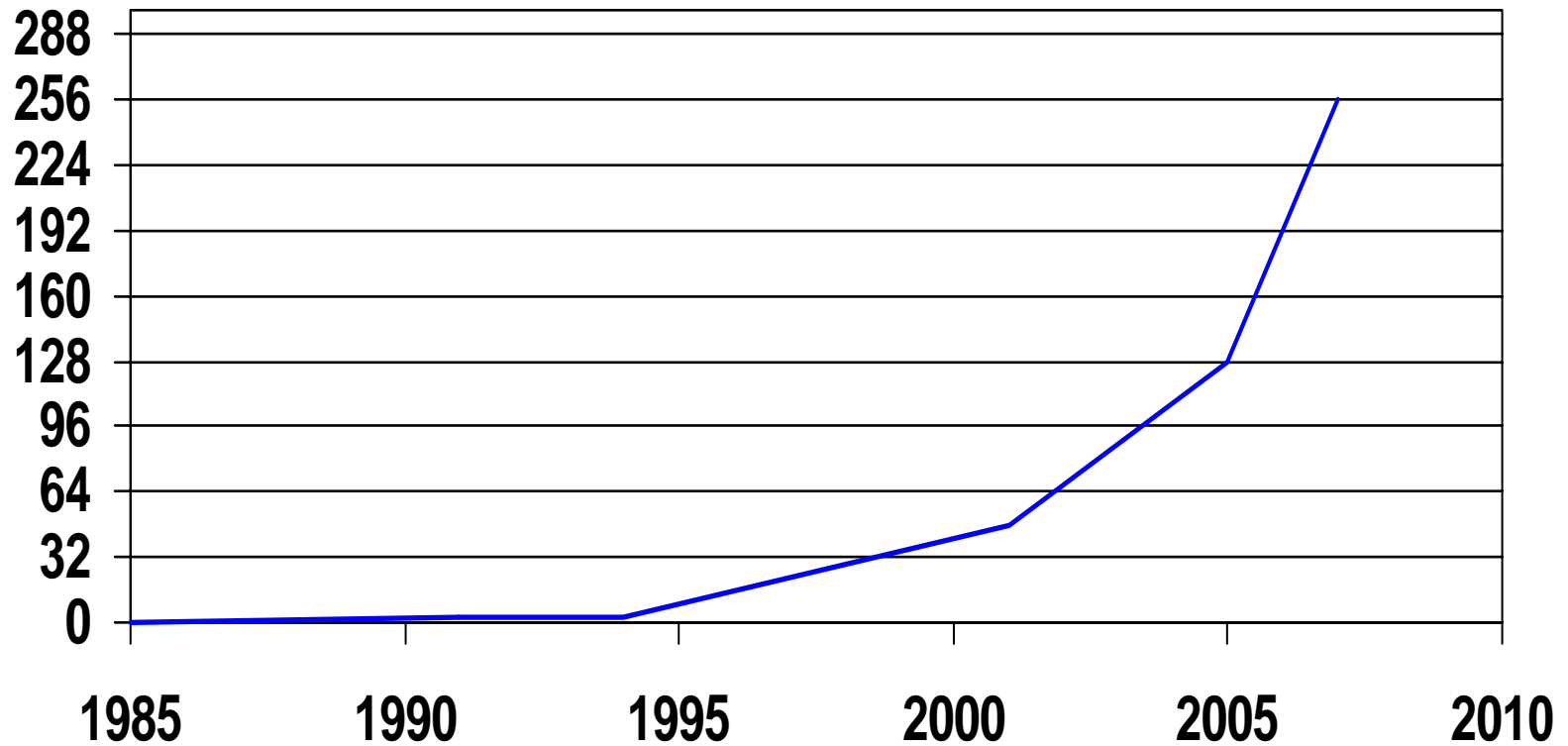
- **Active guest real limit imposed by PTRM space limits (architected): 8 TB**
 - 16 4-GB PTRM spaces; each PTRM space can map 512 GB of guest real
- **Virtual to real ratio (practical): about 3:1**
 - Unless you really, really do your homework on your paging subsystem
 - VMRM-CMM can help with this too (it encourages Linux to Diag x'10' its guest real)

Memory

- **Paging space (architected) (optimal when $\leq 50\%$ allocated):**
 - 11.2 TB for ECKD
 - 15.9 TB for Emulated FBA on FCP SCSI
- **Concurrent paging I/Os per paging volume: 1 for ECKD, >1 for EDEV (Have observed 1.6)**
- **System Execution Space (SXS) (architected): 2 GB**
 - For practical purposes it is 2GB, but there are structures in the space placed above 2GB
- **DCSS aggregate size (architected):**
 - Individual Segments up to 2047 MB
 - Segments can reside above 2GB, starting in z/VM 5.4.0
- **Minidisk Cache (architected): 8GB**

Memory Scaling

Effective Real Memory Use Limits



Page Slots: FCX146 AUXLOG

FCX146 Run 2007/09/06 14:00:28

AUXLOG

Auxiliary Storage Utilization, by Time

From 2007/09/04 09:07:00

To 2007/09/04 10:00:00

For 3180 Secs 00:53:00

Interval	<Page Slots>		<Spool Slots>		<Dump Slots>		<----- Spool Files ----->				<Average MLOAD>	
	Total Slots	Used %	Total Slots	Used %	Total Slots	Used %	<--Created--> Total	/s	<--Purged--> Total	/s	Paging msec	Spooling msec
>>Mean>>	87146k	44	5409096	52	0	..	54	.02	54	.02	2.8	.8
09:08:00	87146k	44	5409096	52	0	..	1	.02	1	.02	2.3	.8
09:09:00	87146k	44	5409096	52	0	..	1	.02	1	.02	3.9	.8
09:10:00	87146k	44	5409096	52	0	..	1	.02	1	.02	3.6	.8
09:11:00	87146k	44	5409096	52	0	..	1	.02	1	.02	2.8	.8
09:12:00	87146k	44	5409096	52	0	..	1	.02	1	.02	2.9	.8

1. This system is using 44% of its page slots.

DASD I/O: FCX109 DEVICE CPOWNED

FCX109 Run 2007/09/06 14:00:28

DEVICE CPOWNED

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Load and Performance of CP Owned Disks

From 2007/09/04 09:07:00

To 2007/09/04 10:00:00

For 3180 Secs 00:53:00

CPU 2094-700

z/VM V.5.3.0 SLU 0701

Page / SP00L Allocation Summary

PAGE slots available	87146k	SP00L slots available	5409096
PAGE slot utilization	44%	SP00L slot utilization	52%
T-Disk cylinders avail.	DUMP slots available	0
T-Disk space utilization	...%	DUMP slot utilization	..%

< Device Descr. ->		<----- Rate/s ----->										User	Serv	MLOAD	Block		
%Used																	
Addr	Devtyp	Volume	Area	Area	Used	<--Page-->		<--Spool-->		Total	SSCH	Inter	Queue	Time	Resp	Page	for
		Serial	Type	Extent	%	P-Rds	P-Wrt	S-Rds	S-Wrt		+RSCH	feres	Lngh	/Page	Time	Size	Alloc
F08B	3390	VS2P49	PAGE	0-3338	45	2.6	1.7	4.4	1.6	1	.02	2.4	2.4	7	89
F090	3390	VS2P69	PAGE	0-3338	45	2.7	1.6	4.3	1.6	1	0	2.7	2.7	7	84

V:R Ratio and Segment Tables: FCX113 UPAGE

FCX113 Run 2007/09/06 14:00:28

UPAGE

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User Paging Activity and Storage Utilization

From 2007/09/04 09:07:00

VS2

To 2007/09/04 10:00:00

CPU 2094-700 SN 2BFBD

For 3180 Secs 00:53:00

z/VM V.5.3.0 SLU 0701

Userid	Data Owned	Paging Activity/s							Number of Pages							Stor Size	Nr of Users	
		<Page Rate>		Page	<---Page Migration-->				<-Resident->		<--Locked-->							
	Spaces	Reads	Write	Steals	>2GB>	X>MS	MS>X	X>DS	WSS	Resrvd	R<2GB	R>2GB	L<2GB	L>2GB	XSTOR	DASD		
>System<	.0	1.7	1.1	4.1	.0	2.4	3.7	1.4	122050	0	2347	106962	6	24	12240	179131	1310M	212
DATAMOVF	.0	.0	.0	.0	.0	.0	.1	.0	13	0	0	0	0	0	483	254	32M	
DATAMOVA	.0	.0	.0	.0	.0	.5	.5	.0	147	0	0	0	0	0	220	368	32M	
DATAMOVB	.0	.0	.0	.0	.0	.6	.6	.0	192	0	0	0	0	0	220	366	32M	
DATAMOV C	.0	.0	.0	.0	.0	.6	.6	.0	191	0	0	0	0	0	220	369	32M	
DATAMOVD	.0	.0	.0	.0	.0	.6	.6	.0	189	0	0	0	0	0	220	362	32M	

1. Resident Guest Pages = (2347 + 106962) * 212 = 88.3 GB
2. V:R = (1310 MB * 212) / 91 GB = 2.98
3. Segment Table Pages: hard to say. Worst case (all 8 GB guests):
212 guests * (4 ST/guest * 4 pg/ST) = 13 MB

PTRM Space: FCX134 DSPACESH

FCX134 Run 2007/09/06 14:00:28

DSPACESH

Shared Data Spaces Paging Activity

From 2007/09/04 09:07:00

To 2007/09/04 10:00:00

For 3180 Secs 00:53:00
0701

CPU 2094-700

z/VM V.5.3.0 SLU

		<----- Rate per Sec. ----->						<-----Number of Pages----->								
Owning								<--Resid-->			<-Locked-->		<-Aliases-->			
userid	Data Space Name	Pgstl	Pgrds	Pgwrt	X-rds	X-wrt	X-mig	Total	Resid	R<2GB	Lock	L<2GB	Count	Lockd	XSTOR	DASD
>System<	-----	.026	.016	.001	.015	.026	.000	103k	1208	51	0	0	0	0	34	4981
SYSTEM	FULL\$TRACK\$CACHE\$1	.000	.000	.000	.000	.000	.000	524k	0	0	0	0	0	0	0	0
SYSTEM	ISFCDATASPACE	.000	.000	.000	.000	.000	.000	524k	113	8	8	8	113	100	0	27
SYSTEM	PTRM0000	4.257	.492	.442	3.957	4.036	.000	1049k	386k	15885	0	0	0	0	5195	683k
SYSTEM	REAL	.000	.000	.000	.000	.000	.000	24M	0	0	0	0	0	0	0	0
SYSTEM	SYSTEM	.080	.001	.034	.079	.080	.000	524k	45	10	0	0	44	0	47	510k

1. PTRM space = (386,000 + 15885) = 401,885 = 1.53 GB
(NB: this is z/VM 5.3)

Real Memory: FCX254 AVAILLOG

FCX254 Run 2007/09/06 14:00:28

AVAILLOG

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Available List Management, by Time

From 2007/09/04 09:07:00

To 2007/09/04 10:00:00

For 3180 Secs 00:53:00

CPU 2094-700

z/VM V.5.3.0 SLU 0701

----- Available List Management -----																				
----- Thresholds -----				----- Page Frames -----						----- Times -----		----- Replenishment -----						----- Perct -----		
Interval	<---Low---	<---Hi gh---		<Avai lable>	<Obtai ns/s>	<Returns/s>				<-Empty->		<---Scan1---	<---Scan2---	<-Em-Scan-->	Scan	Emerg				
>>Mean>>	<2GB	>2GB	<2GB	>2GB	<2GB	>2GB	<2GB	>2GB	<2GB	>2GB	<2GB	>2GB	Compl	Pages	Compl	Pages	Compl	Pages	Fai l	Scan
>>Mean>>	20	7588	5820	13388	5130	7678	323.3	857.4	311.5	844.8	0	0	27	1381k	63	1380k	58	84490	82	88
09:08:00	20	7680	5820	13480	6665	15122	353.3	838.5	353.2	1007	0	0	0	43091	3	26491	0	0	3	100
09:09:00	20	7680	5820	13480	3986	5496	163.1	640.2	108.9	442.7	0	0	1	14528	0	0	0	0	0	0
09:10:00	20	7681	5820	13481	6622	9542	222.4	556.1	257.0	598.3	0	0	0	30103	2	8868	0	0	1	100
09:11:00	20	7681	5820	13481	4982	6710	292.1	615.2	248.8	533.6	0	0	0	21246	0	8547	1	3989	1	100
09:12:00	20	7681	5820	13481	4769	1560	284.9	946.9	254.4	830.0	0	0	0	18253	0	22438	2	656	1	100

1. Pct ES = 88% generally this system is tight on storage
2. Scan fail >0 generally this system is tight on storage
3. Times Empty = 0 this indicates it isn't critical yet

SXS Space: FCX261 SXS AVAIL

FCX261 Run 2007/09/06 14:00:28

SXS AVAIL

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System Execution Space Page Queues Management

From 2007/09/04 09:07:00

To 2007/09/04 10:00:00

For 3180 Secs 00:53:00

CPU 2094-700

z/VM V.5.3.0 SLU 0701

Interval	<-- Backed <2GB Page Queue -->					<-- Backed >2GB Page Queue -->					<----- Unbacked Page Queue ----->								
	Avail	<-Pages/s-->	<Preferred>	Pages Taken	Return	Avail	<-Pages/s-->	<Preferred>	Pages Taken	Return	Pages	<-Pages/s-->	<Preferred>	Used	Empty	Thres	Att/s	Stolen	MinPgs
>>Mean>>	26	.513	.509	.513	.000	3	1.798	1.804	1.798	4.114	466946	130.3	130.1	126.2	.000	128	.000	128	...
09:08:00	26	.483	.383	.483	.000	0	1.650	1.650	1.650	3.667	467829	128.2	127.3	124.5	.000	128	.000	128	...
09:09:00	26	.500	.500	.500	.000	0	.583	.583	.583	3.067	465679	120.8	84.98	117.8	.000	128	.000	128	...
09:10:00	27	.517	.533	.517	.000	0	1.183	1.183	1.183	4.000	467657	109.1	142.1	105.1	.000	128	.000	128	...
09:11:00	27	.517	.517	.517	.000	0	1.633	1.633	1.633	2.917	467632	137.2	136.8	134.3	.000	128	.000	128	...
09:12:00	29	.450	.483	.450	.000	0	2.000	2.000	2.000	3.383	467654	129.9	130.2	126.5	.000	128	.000	128	...
09:13:00	27	.517	.483	.517	.000	0	2.483	2.483	2.483	3.550	467698	139.3	140.0	135.7	.000	128	.000	128	...
09:14:00	25	.550	.517	.550	.000	0	2.000	2.000	2.000	2.750	465651	119.0	84.92	116.3	.000	128	.000	128	...

1. How we touch guest pages: (1) 64-bit; (2) AR mode; (3) SXS.
2. There are 524,288 pages in the SXS.
3. This system has 466,000 SXS pages available on average.

FCX178 Run 2008/04/15 10:00:22 MDCSTOR Page 76

Mini disk Cache Storage Usage, by Time

From 2008/04/15 09:47:11

To 2008/04/15 10:00:11

CPU 2084-320 SN 17F2A

For 780 Secs 00:13:00

z/VM V.5.3.0 SLU 0000

<----- Main Storage Frames ----->

Interval	<--Actual-->			Min	Max	Page	Steal	
End Time	Ideal	<2GB	>2GB	Set	Set	Del /s	Invokd/s	Bias
>>Mean>>	5839k	82738	1354k	0	7864k	0	.000	1.00
09:57:41	5838k	119813	1932k	0	7864k	0	.000	1.00
09:58:11	5838k	119813	1932k	0	7864k	0	.000	1.00
09:58:41	5838k	119825	1932k	0	7864k	0	.000	1.00
09:59:11	5838k	119825	1932k	0	7864k	0	.000	1.00
09:59:41	5838k	119825	1932k	0	7864k	0	.000	1.00
10:00:11	5838k	119837	1932k	0	7864k	0	.000	1.00

- Xstore not used for this configuration so edited out from report.
- Add up the pages in Main Storage and you get ~8GB

FCX134 Run 2008/04/15 10: 00: 22

DSPACESH

Shared Data Spaces Paging Acti vi ty

From 2008/04/15 09: 47: 11

To 2008/04/15 10: 00: 11

For 780 Secs 00: 13: 00

This is a performance report for system XYZ

		-----Number of Pages----->										
Owni ng		Users	<--Resid-->				<-Locked-->		<-Al i ases-->			
Userid	Data Space Name	Permt	Total	Resid	R<2GB	Lock	L<2GB	Count	Lockd	XSTOR	DASD	
>System<	-----	0	1507k	5665	101	0	0	100	0	0	0	
SYSTEM	FULL\$TRACK\$CACHE\$1	0	524k	0	0	0	0	0	0	0	0	
SYSTEM	FULL\$TRACK\$CACHE\$2	0	524k	0	0	0	0	0	0	0	0	
SYSTEM	FULL\$TRACK\$CACHE\$3	0	524k	0	0	0	0	0	0	0	0	
SYSTEM	FULL\$TRACK\$CACHE\$4	0	524k	0	0	0	0	0	0	0	0	
SYSTEM	ISFCDATASPACE	0	524k	0	0	0	0	0	0	0	0	
SYSTEM	PTRM0000	0	1049k	44489	0	0	0	0	0	0	0	
SYSTEM	REAL	0	7864k	0	0	0	0	0	0	0	0	
SYSTEM	SYSTEM	0	524k	805	787	0	0	800	0	0	0	
SYSTEM	VIRTUAL\$FREE\$STORAGE	0	524k	23	23	0	0	0	0	0	0	

- You'll see the address spaces used for MDC (track cache)
- Values here are zero for page counts, ignore.

Reorder Processing - Background

- **Page reorder** is the process in z/VM of managing user frame owned lists as input to demand scan processing.
 - It includes resetting the HW reference bit.
 - Serializes the virtual machine (all virtual processors).
 - In all releases of z/VM
- **It is done periodically on a virtual machine basis.**
- **The cost of reorder is proportional to the number of resident frames for the virtual machine.**
 - Roughly 130 ms/GB resident
 - Delays of ~1 second for guest having 8 GB resident
 - This can vary for different reasons +/- 40%

Reorder Processing - Diagnosing

■ Performance Toolkit

- Check resident page fields (“R<2GB” & “R>2GB”) on FCX113 UPAGE report
 - Remember, Reorder works against the resident pages, not total virtual machine size.
- Check Console Function Mode Wait (“%CFW”) on FCX114 USTAT report
 - A virtual machine may be brought through console function mode to serialize Reorder. There are other ways to serialize for Reorder and there are other reasons that for CFW, so this is not conclusive.

■ REORDMON

- Available from Bill Bitner now and the VM Download Page <http://www.vm.ibm.com/download/packages/> in near future
- Works against raw MONWRITE data for all monitored virtual machines
- Works in real time for a specific virtual machine
- Provides how often Reorder processing occurs in each monitor interval

Reorder Processing - Mitigations

- **Try to keep the virtual machine as small as possible.**
- **Virtual machines with multiple applications may need to be split into multiple virtual machines with fewer applications.**
- **Known requirement at IBM to bring relief in this area.**

I/O

- **Number of subchannels in a partition (aka device numbers) (architected): 65,535**
- **CHPIDs per server (z9 EC): 336 FICON or 1024 ESCON**
- **Device numbers per disk volume**
 - Without PAV, 1
 - With PAV or HyperPAV, 8 (base plus seven aliases)
- **Virtual Devices per Virtual Machine:**
 - 24576 (24K)
- **Concurrent real I/Os per ECKD disk volume: 1 usually, but 8 with PAV or HyperPAV if of guest origin**
- **Concurrent real I/Os per chpid (aka “open exchange limit”)**
 - 1 for ESCON
 - 64 for FICON
- **I/O rates:**
 - Fastest FICON is a 4 Gb link
 - About 1-2 msec per I/O are required for a nominal DASD I/O from a z9 to a 2107 (500-1000/sec/device)

I/O

▪ Volume sizes

- Largest ECKD minidisk that can contain a CMS file system (architected): 32768 cylinders (22.5 GB)
- Largest EFBA minidisk that can contain a CMS file system (architected): 381 GB
 - Practical limit of 22GB due to file system structure under 16MB, unless there are very few files.
- Largest ECKD volume, period: 65536 cylinders (43 GB)
- Largest EDEV CP can use: 1024 GB (but PAGE, SPOL, DRCT must be below 64 GB line on volume)
- Largest EDEV, period: 2^{32} FB-512 blocks (2048 GB)
- VDISK size (architected): 2 GB (minus eight 512-byte blocks)

▪ Total VDISK (architected): 2TB

▪ Single VSWITCH OSAs: 8

DASD I/O: FCX108 DEVICE

FCX108 Run 2007/09/06 14:00:28

DEVICE

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General I/O Device Load and Performance

From 2007/09/04 09:07:00

To 2007/09/04 10:00:00

CPU 2094-700 SN

For 3181 Secs 00:53:01

z/VM V.5.3.0 SLU 0701

<-- Device Descr. -->			Mdisk Pa-	<-Rate/s->				<----- Time (msec) ----->				Req.	<Percent>		SEEK	Recov	<-Throttle->		
Addr	Type	Label /ID	Links	ths	I/O	Avoid	Pend	Disc	Conn	Serv	Resp	CUWt	Qued	Busy	READ	Cyls	SSCH	Set/s	Dly/s
>>	All	DASD <<5	.4	.2	.1	3.4	3.7	3.7	.0	.0	0	17	1173	00
F024	3390	VS2426	1	4	12.9	147.0	.2	.7	.4	1.3	1.3	.0	.0	2	91	193	0
OC20	CTCA		...	1	12.63	.2	.6	1.1	1.1	.0	.0	1	0
F685	3390	VS2W01	290	4	11.8	.3	.2	.0	.3	.5	.5	.0	.0	1	84	89	0
F411	3390	VS2613	1	4	10.6	.5	.2	.3	.4	.9	.9	.0	.0	1	1	1303	0

Other

- **Number of spool files (architected):**
 - 9999 per user
 - 1.6 million spool files per system
 - 1024 files per warm start block * (180 * 9) warm start blocks
- **Number of logged-on virtual machines (approximate): about 100,000 (per designers)**

Metrics for Formal Spin Locks

FCX265 CPU 2094 SER 19B9E Interval 02:31:51 - 12:34:01 GDLVM7

```

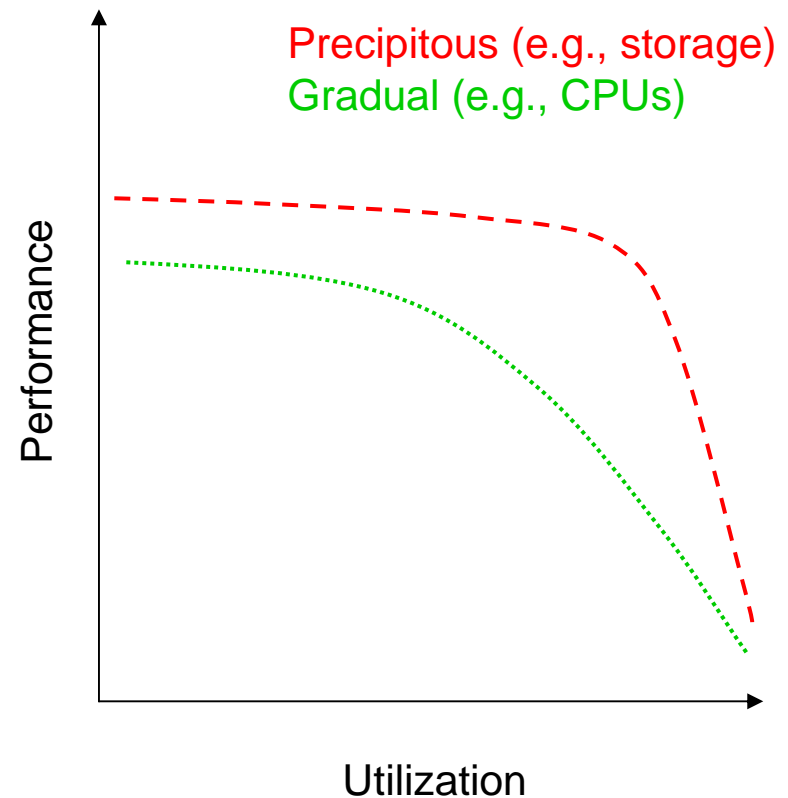
<----- Spin Lock Activity ----->
<----- Total -----> <--- Exclusive ---> <----- Shared ----->
Interval                Locks Average   Pct  Locks Average   Pct  Locks Average   Pct
End Time LockName      /sec   usec  Spin  /sec   usec  Spin  /sec   usec  Spin
>>Mean>> SRMATDLK      1.9    .539  .000  1.9    .539  .000   .0    .000  .000
>>Mean>> RSAAVCLK       .0    2.015  .000   .0    2.015  .000   .0    .000  .000
>>Mean>> FSDVMLK        .0   24.97  .000   .0   24.97  .000   .0    .000  .000
>>Mean>> SRMALOCK       .0     .000  .000   .0     .000  .000   .0    .000  .000
>>Mean>> HCPTRQLK       4.1     .195  .000   4.1     .195  .000   .0    .000  .000
>>Mean>> SRMSLOCK      34.0    1.096  .001  32.7    1.037  .001   1.3    .001  .000
    
```

Latent Limits

- **Sometimes it's not an architected limit**
- **Sometimes it's just “your workload won't scale past here, because...”**
- **In our studies of z/VM 5.3, we found these kinds of latent limits:**
 - Searching for a below-2-GB frame in lists dominated by above-2-GB frames (storage balancing functions)
 - Contention for locks, usually the scheduler lock
- **These kinds of phenomena were the reasons we published the limits to be 256 GB and 32 engines**
 - We wanted to publish supported limits we felt would be safe in a very large variety of workloads and environments
 - Many of our measurement workloads scaled higher than this (for example, 440 GB and 54 engines)

Other Notes on z/VM Limits

- **Sheer hardware:**
 - z/VM 5.2: 24 engines, 128 GB real
 - z/VM 5.3: 32 engines, 256 GB real
 - zSeries: 65,000 I/O devices
- **Workloads we've run in test have included:**
 - 54 engines
 - 440 GB real storage
 - 128 GB XSTORE
 - 240 1-GB Linux guests
 - 8 1-TB guests
- **Utilizations we routinely see in customer environments**
 - 85% to 95% CPU utilization without worry
 - Tens of thousands of pages per second without worry
- **Our limits tend to have two distinct shapes**
 - Performance drops off slowly with utilization (CPUs)
 - Performance drops off rapidly when wall is hit (storage)



Keeping Tabs on Consumption Limits

- **Processor**

- CPU utilization: FCX100 CPU, FCX114 USTAT

- **Memory & Paging**

- Page slots in use: FCX146 AUXLOG
- DASD I/O: FCX109 DEVICE CPOWNED
- V:R Memory ratio: FCX113 UPAGE
- PTRM space consumed: FCX134 DSPACESH
- Storage in use for segment tables: FCX113 UPAGE
- Consumption of SXS space: FCX261 SXS AVAIL
- MDC: FCX178 MDCSTOR, FCX134 DSPACESH
- Consumption of real memory: FCX103 STORAGE, FCX254 AVAILLOG
- Consumption of expanded storage: FCX103 STORAGE

- **I/O**

- DASD I/O: FCX108 DEVICE
- Concurrency on FICON chpids: FCX131 DEVCONF, FCX215 INTERIM FCHANNEL, FCX168 DEVLOG

What Consumption Limits Will We Hit First?

- **Depends on workload**
 - Guest-storage-intensive:
 - page slots on DASD... at 5-6 TB things start to get interesting... mitigate by paging to SCSI
 - utilization on paging volumes and chpids: watch for MLOAD elongation... mitigate by spreading I/O
 - mitigation by application tuning... perhaps smaller guests
 - segment table constraints: probably an issue at 50% (64 TB of logged-on guest real) ... not anytime soon
 - Real-storage-intensive:
 - Ability of the system to page will limit you: ensure adequate XSTORE and paging capacity
 - You can define > 256 GB of real storage, but we are aware that some workloads cannot scale that high
 - Mitigation by application tuning or by using CMM
 - CPU-intensive:
 - FCX100 CPU and FCX 114 USTAT will reveal CPU limitations
 - You can define > 32 engines, but we are aware that some workloads cannot scale that high
 - Mitigation by application tuning
 - I/O-intensive:
 - Device queueing: consider whether PAV or HyperPAV might offer leverage
 - Chpid utilization: add more chpids per storage controller
 - Ultimately partitions can be split, but we would prefer you not have to do this (too complicated)
- **Without trend data (repeated samples) for *your* workloads it is difficult to predict which of these limits *you* will hit first**

Summary

- **Knowing Limits:**
 - Real resource consumption
 - Limits to managing the virtualization of real resources
- **Measuring Limits:**
 - Knowing where to watch for these limits
 - Including these in capacity planning
- **Managing Limits**
 - Tuning and configuring
 - Planning for growth