

## z/VM Tuning Revisited with Specialty Engines for z/OS Session 9122

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8/20/2009



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# Agenda

- Some general z/VM Scheduling and Dispatching Discussion
- Some general z/OS Guest Tuning Discussion
- Background on Specialty Engine support in z/VM
- Tuning of Specialty Engines
- Miscellaneous z/OS Tuning



# z/VM Scheduling & Dispatching at the High Level

### Objectives of the z/VM Scheduler

- Protect the system from over committing resources to where the system thrashes
- Prioritize access to system resources

### Objectives of the z/VM Dispatcher

- Effectively run virtual processors based on their priorities



## **Scheduler Protection from Thrashing**





# Three Main Controls to Entering Dispatch List

cp q srm

- IABIAS : INTENSITY=90%; DURATION=2
- LDUBUF : Q1=100% Q2=75% Q3=60%
- STORBUF: Q1=125% Q2=105% Q3=95%
- DSPBUF : Q1=32767 Q2=32767 Q3=32767

DISPATCHING MINOR TIMESLICE = 5 MS

- MAXWSS : LIMIT=9999%
- ..... : PAGES=999999
- XSTORE : 0%

Ready;

LDUBUF: protects from thrashing DASD Paging

STORBUF: protects from general thrashing of real memory

DSPBUF: Absolute number allowed in dispatch list for each scheduling class



# Comments on SRM Value for z/OS Systems

- Defaults were determined based on traditional workload with mix of interactive CMS and Guest work.
- Potential benefit from changing SRM values.
  - If having problems, investigate STORBUF first
  - Second, look at LDUBUF
  - Keep your hands off DSPBUF unless you really know what you are doing.
  - Avoid temptation to increase/change several values at once
- Increasing DSPSPLICE was considered clever at one time. The overhead from dispatching these days probably isn't worth the downside of increasing it. Leave it alone.



# Deadline Scheduling – Prioritizing Work

- Each virtual processor has a priority computed as a 'deadline' for when a unit of work should be completed.
- This 'deadline' is a time value on an artificial TOD often referred to as ATOD
- The 'deadline' is computed based on several factors, but the most significant is the normalized Share value
- Therefore the share setting is a big knob
- Virtual processors get ordered for dispatching based on their deadlines



# ATOD and Deadline

### ATOD



Simplified offset formula used to set deadline 'offset' from current ATOD:

```
OFFSET = Minor_TimeSlice + Previous_TimeSlice Overrun
```

Normalized\_Share x Number\_PUs



# A Word About QUICKDSP

- Quick Dispatch (SET QUICKDSP) for a virtual machine allows it to pass from eligible list to dispatch list without going through the system resource checks.
- Does NOT turn off the scheduler completely.
- Should be set on for:
  - Mission Critical Virtual Machines
  - Virtual Machines that are extensions of Operating System (e.g. RACF, TCP/IP)
  - A virtual machine you have access to for tuning and problem determination.



#### z/VM and Specialty Engine Support



8/20/2009



# z/VM 5.3.0 Support for Specialty Processors

- z/VM V5.3 introduces support for zAAP and zIIP specialty processors
  - System z Application Assist Processors (zAAPs) provide an economical Java execution environment for z/OS and z/OS.e
  - System z9 Integrated Information Processors (zIIPs) designed to help improve resource optimization and lower the cost for eligible z/OS and z/OS.e workloads by offloading software system overhead from standard Central Processors (CPs); this includes certain DB2 processing

#### z/VM support is provided for z/OS guest exploitation

- Offers additional hardware support for z/OS-on-z/VM development and test support

#### Two levels of z/VM support:

- Simulation support
  - z/VM dispatches virtual zAAPs and zIIPs on real CP engines
  - Only possible if the underlying hardware is capable of supporting the real engine type
  - Does not require activation of real specialty engine(s) within the mainframe server
- Virtualization support
  - z/VM dispatches virtual zAAPs and zIIPs on corresponding real specialty engines
- Consistent with z/OS, there are no z/VM license fees associated with real zAAP or zIIP processors



### z/VM 5.3.0 Specialty Processor Support Example Simulating Specialty Engines in Virtual Machines

- Allows users to discover the operational aspects of using zAAPs and zIIPs in a z/OS environment without having to purchase real specialty processors
- May help users assess specialty-processor eligible workloads in a z/OS environment
- Provides a function test environment for z/OS workloads that use specialty processors
- Consumes CP processor capacity to host virtual zAAP and zIIP processor cycles





### z/VM 5.3.0 Specialty Processor Support Example Using Real Specialty Engines in Virtual Machines

- Allows users to test and verify z/OS specialty processor support on the real hardware
- Users can maximize real specialty processor utilization by sharing processors among production and test LPARs
- Consumes specialty processor capacity to host virtual zAAP and zIIP processor cycles





# z/VM-Mode LPAR Support for IBM System z10

New LPAR type for IBM System z10: *z/VM-mode* 

Allows z/VM V5.4 users to configure all CPU types in a z10 LPAR

#### Offers added flexibility for hosting mainframe workloads

- Add IFLs to an existing standard-engine z/VM LPAR to host Linux workloads
- Add CPs to an existing IFL z/VM LPAR to host z/OS, z/VSE, or traditional CMS workloads
- Add zAAPs and zIIPs to host eligible z/OS specialty-engine processing
- Test integrated Linux and z/OS solutions in the same LPAR

#### No change to software licensing

Software continues to be licensed according to CPU type





## Some Additional Background

#### CPU Affinity

- Setting to control whether virtualized Specialty Engines must be dispatched on real processors of that type
- ON means virtual type must equal real type
- Suppressed: you have asked for ON, but we don't have processors of that type to use

#### Processor Type Pools

- Scheduling is done within a pool for CPUAFFINITY ON
- Capacity Planning of each type
- ATOD, ATOD2, etc. for each Processor Type Pool

#### Primary vs. Secondary Processor

- Primary: CP or IFL
- Secondary: zAAP, zIIP, and sometimes IFL (secondary to CPs)

#### Different Speed Processors

- Specialty engines are full-speed on all z9 and z10 machines, while some general purpose run at a fractional speed.
- The z/VM Scheduler is a deadline scheduler, not a consumption scheduler



# Considerations for z/VM-mode LPARs

- Merging IFL only and CP only partitions in a z/VMmode partition requires planning
  - First step, make virtual machines on IFL LPAR have virtual IFLs
  - For duplicated work (RACF, TCP/IP, etc.), need to determine which to use or in some cases which to duplicate
  - Remember that in some environments, the IFLs may be faster than the CPs.
  - Determine any changes you want to make to the charge back model.

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## **Output from INDICATE USER EXPanded**

CPU 00: Ctime=0 00:00:53 Vtime=0 00:00:00 Ttime=0 00:00:00 Rdr=0 Prt=0 Pch=0 IO=332 Type=CP CPUAffinity=ON VtimePrimary=0 00:00:00 TtimePrimary=0 00:00:00 VtimeSecondary=0 00:00:00 TtimeSecondary=0 00:00:00 CPU 01: Ctime=0 00:00:30 Vtime=0 00:00:00 Ttime=0 00:00:00 Rdr=0 Prt=0 Pch=0 IO=0 Type=ZAAP CPUAffinity=SUPP VtimePrimary=0 00:00:00 TtimePrimary=0 00:00:00 VtimeSecondary=0 00:00:00 TtimeSecondary=0 00:00:00 CPU 02: Ctime=0 00:00:20 Vtime=0 00:00:00 Ttime=0 00:00:00 Rdr=0 Prt=0 Pch=0 IO=0 Type=ZIIP CPUAffinity=SUPP VtimePrimary=0 00:00:00 TtimePrimary=0 00:00:00 VtimeSecondary=0 00:00:00 TtimeSecondary=0 00:00:00



# **Specialty Engines and Share Settings**

- The Share setting for a virtual machine applies to each pool of the processor types
  - CP, IFL, zIIP, zAAP, etc.
- z/VM 5.3.0 one share setting for all types
- z/VM 5.4.0 added support to set a separate share setting for each processor type pool
  - Default is TYPE ALL and results in one setting for all types, effectively like z/VM 5.3.0
- Normalized to the sum of shares of virtual machines in dispatch list for each pool of the processor types
- Absolute (and normalized) is percentage of resources of a given processor type.



NN% = (IPW) In Perfect World percentage of real processor

NN% = split of share per virtual processor







NN% = IPW percentage of real processor NN% = split of share per virtual processor







#### NN% = IPW percentage of real processor NN% = split of share per virtual processor







NN% = IPW percentage of real processor NN% = split of share per virtual processor





NN% = IPW percentage of real processor NN% = split of share per virtual processor NN = relative split of share per virtual processor





#### NN% = IPW percentage of real processor NN% = split of share per virtual processor NN = relative split of share per virtual processor



Sum of CP Relative Shares = 800 Sum of zAAP Relative Shares = 300





#### NN% = IPW percentage of real processor NN% = split of share per virtual processor NN = relative split of share per virtual processor



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Total Processor for SERVERB is 15+15+15+15+20 = 80%

Total Processor for SERVERC is 21+21+21 = 63%









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## FCX126 LPAR – updated

FCX126 Run 2007/06/20 09:55:12					]	LPAR																	
														Logi	cal Pa	Partition Activity							
Parti	tion Nr	·. τ	Jpid	#Proc	Weight	Wait-C	Cap	%Load	CPU	%Busy	%0vhd	%Susp	%VMld	%Logld	Туре								
KST1		4	04	5	20	NO	NO	51.5	0	99.9	.0	.1	99.9	99.9	CP								
					20		NO		1	99.9	.0	.1	99.8	99.9	CP								
					20		NO		2	99.9	.1	.1	99.8	99.9	CP								
					80		NO		3	.0	.0	.4	.0	.0	ZAAP								
					80		NO		4	9.3	.2	.3	9.0	9.0	ZIIP								
Summa	ry of p	hysi	cal	proces	ssors:																		
Туре	Number	We	eight	Ded:	icated			•	New	Туре	e for	each	n pro	cessoi	2								
CP	3	5	100	)	0		Totals for Processor type																
ZAAP	1		100	)	0				1000				,DOT	eypeb									
IFL	1		0	)	0				One	IFL	on m	achin	ie, bi	ut no									
ZIIP	1		100	)	0		active partitions using it.							•									

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## FCX202 LPARLOG – updated

FCX202 Run 2007/06/20 09:55:12

LPARLOG

Logical Partition Activity Log

Interval	terval <partition-></partition->							<- Load per Log. Processor>						
End Time	Name	Nr.	Upid	#Proc	Weight	Wait-C	Cap	%Load	%Busy	%0vhd	%Susp	%VMld	%Logld	Туре
>>Mean>>	KST1	4	04	5	80	NO	NO	•••	61.8	.1	.2	61.7	61.7	MIX
>>Mean>>	Total	••	••	6	100	••	••	.1	30.9	.0	• • •	• • •	• • •	••



## FCX180 SYSCONF

FCX180 Run 2007/06/20 09:53:49

SYSCONF

System Configuration

Initial Status on 2007/03/07 at 21:30, Processor 2096-X03 Real Proc: Cap 2224, Total 7, Conf 3, Stby 0, Resvd 4 Sec. Proc: Cap 1760, Total 3, Conf 3, Stby 0, Resvd 2

- With new Specialty Engine support, your z/VM system may include processors that are different speeds.
- Smaller "Cap" number indicates faster Processor.



## FCX144 PROCLOG – updated

FCX144 Run 2007/06/20 09:53:49

Drogoggor	Nativity	hu	Timo
Processor	ACTIVITY,	DY	Time

PROCLOG

< Percent Busy> < Rates per Sec>									••>	
	C									
Interval	Р						Inst			
End Time	U Type	Total	User	Syst	Emul	Vect	Siml	DIAG	SIGP	SSCH
>>Mean>>	0 CP	99.8	99.5	.3	97.9	••••	125.0	12.6	.7	71.6
>>Mean>>	1 CP	99.8	99.5	.2	98.0	••••	120.9	4.5	.8	58.4
>>Mean>>	2 CP	99.8	99.5	.3	98.0	••••	123.4	3.2	.7	59.5
>>Mean>>	3 ZAAP	96.0	96.0	.1	95.8	• • • •	1.1	.0	36.6	1.4
>>Mean>>	4 ZIIP	8.8	8.4	.4	8.1	••••	1.0	.0	289.9	7.5
>>Mean>>	. CP	99.7	99.5	.2	98.0	••••	123.0	6.7	.7	63.1
>>Mean>>	. ZAAP	96.0	96.0	.1	95.8	••••	1.1	.0	36.6	1.4
>>Mean>>	. ZIIP	8.8	8.4	.4	8.1	• • • •	1.0	.0	289.9	7.5



## Virtual Sysplex Environments

- Key is tuning effectively for the virtual coupling machines
  - QUICKDSP ON
  - Sufficiently High Share setting
  - Using real ICFs in z/VM-mode LPARs in z/VM 5.4.0 where appropriate
- Beware of scenarios with both a large number of systems in a virtual sysplex and the systems join and leave the sysplex frequently.
  - High CP CPU overhead as the z/OS systems that are not changing state issue large number of messages while the coupling machine is busy making updates for the system that is leaving/joining.
  - Privileged Operations count will be very high (>10,000s/second)



## **Miscellaneous Tuning Thoughts**

### Disable IRD for z/OS virtual machines

#### PAV Usage & Trade-offs

- Dedicating volumes to z/OS guests and letting them use PAV can be the best performance
- More flexibility in sharing volumes and using PAV volumes through the minidisk support
- Analysis of where the I/Os are queued up may require looking at both z/VM data and z/OS data

### HiperDispatch

Does not apply in z/VM guests



# Summary

- Specialty Engines enhance z/VM's virtualization capabilities
- A few things to keep in mind...
  - Share for virtual machine applies to each processor type pool
  - There are scenarios where processors can be different speeds
  - Looking at averages takes on new meaning
  - CPU Affinity Setting is important
- Monitor and Accounting records updated to provide needed information
- For more on scheduling, see VM Scheduler Made Simple
  - http://www.vm.ibm.com/devpages/bitner/presentations/vmsched.html
- For more on Specialty Engine Performance, see Performance Report
  - http://www.vm.ibm.com/perf/reports/zvm/html/530se.html